

## NEW CONSTRUCTIONS OF FUNCTION AND SET CHAINS

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**Abstract:** Function and set chains can be constructed using ordinal notation systems. Here, earlier results are improved, and longer chains constructed. In particular, a simpler method is given for constructing set chains.

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### 1. Function Chains for Schemes

Function chains for schemes have been given in [4]. These will be reviewed here, with some changes. Recall from [6] that a scheme over  $\kappa$  where  $\kappa \in \text{Inac}$  is a pair  $\sigma = \langle \sigma, \phi \rangle$  where  $\sigma < \kappa^+$  and  $\phi$  is a function whose domain is the set of limit ordinals  $\alpha \leq \sigma$ . For  $\alpha \in \text{Dom}(\phi)$ ,  $\phi(\alpha)$  is an increasing function with domain an ordinal  $\eta \leq \kappa$ , and whose range is an unbounded subset of  $\alpha$ . If  $\text{Cf}(\alpha) < \kappa$  then  $\eta < \kappa$ , and if  $\text{Cf}(\alpha) = \kappa$  then  $\eta = \kappa$ .

Recursions and inductions on schemes are broken into four cases, as follows.

Case 0:  $\sigma = 0$ .

Case 1:  $\sigma = \tau + 1$ .  $\tau$  denotes  $\sigma_{\leq \tau}$ .

Case 2:  $\sigma \in \text{Lim}$ ,  $\text{Cf}(\sigma) < \kappa$ . For  $\xi < \eta$   $\sigma_\xi$  denotes  $\phi(\sigma)(\xi)$ ; and  $\sigma_\xi$  denotes  $\sigma_{\leq \sigma_\xi}$ .

Case 3:  $\sigma \in \text{Lim}$ ,  $\text{Cf}(\sigma) = \kappa$ .  $\sigma_\xi$  and  $\sigma_\xi$  are as in case 2, for  $\xi < \kappa$ .

The thin subset  $T_\sigma \subseteq \kappa$  is defined recursively as follows.

- 0:  $\emptyset$ .
- 1:  $T_\tau$ .
- 2:  $(\eta + 1) \cup \cup_{\xi < \eta} T_{\sigma_\xi}$ .
- 3:  $\nabla_{\xi < \kappa} T_{\sigma_\xi}$ .

A function  $f_\sigma : \text{In}_\kappa \mapsto \kappa$  is defined recursively as follows.

- 0: The identically 0 function.
- 1:  $f_\tau + 1$ .
- 2:  $\sup_{\xi < \eta} f_{\sigma_\xi}$ , with  $f_\sigma(\lambda)$  set to 0 if  $\lambda \leq \theta$ .
- 3:  $\text{dsup}_{\xi < \kappa} f_{\sigma_\xi}$ .

This definition was given in [4], for domain  $\kappa$ , and it was observed that these functions are the well-known canonical functions, which are unique mod the thin ideal. As noted in [7], the domain  $\text{In}_\kappa$  is convenient for defining  $f_{\kappa^+}$ . This results in the triviality that when  $\lambda$  is the smallest inaccessible,  $f_\sigma$  is the empty function for all  $\sigma$ . The proviso in case 2 ensures that  $f_\sigma(\lambda) < \lambda^+$ , as a simple induction shows; this does not seem to be necessary, though.

**Lemma 1.** *Suppose  $\sigma', \sigma$  are schemes, and  $\lambda \notin T_{\sigma'} \cup T_\sigma$ .*

- a. *If  $\sigma' \leq \sigma$  then  $f_{\sigma'}(\lambda) \leq f_\sigma(\lambda)$ .*
- b. *If  $\sigma' = \sigma$  then  $f_{\sigma'}(\lambda) = f_\sigma(\lambda)$ .*
- c. *If  $\sigma = \tau + 1$  then  $f_\sigma(\lambda) = f_\tau(\lambda) + 1$ .*
- d. *If  $\sigma' < \sigma$  then  $f_{\sigma'}(\lambda) < f_\sigma(\lambda)$ .*

*Proof.* Part a follows by induction on  $\sigma$ . In case 0,  $\sigma$  is 0 also and the claim is trivial. In case 1, if  $\sigma' \leq \tau$  then inductively  $f_{\sigma'}(\lambda) \leq f_\tau(\lambda)$ , and clearly  $f_\tau(\lambda) < f_\sigma(\lambda)$ . If  $\sigma' = \sigma$  then inductively  $f_{\sigma'}(\lambda) = f_\tau(\lambda)$ , and  $f_{\sigma'}(\lambda) = f_\sigma(\lambda)$  follows. In case 2, if  $\sigma' < \sigma$  then  $\sigma' < \sigma_\xi$  for some  $\xi < \eta$ . Inductively,  $f_{\sigma'}(\lambda) \leq f_{\sigma_\xi}(\lambda)$ , and clearly  $f_{\sigma_\xi}(\lambda) \leq f_\sigma(\lambda)$ . If  $\sigma' = \sigma$ , by the argument just given,  $f_{\sigma'_\xi}(\lambda) \leq f_\sigma(\lambda)$  for all  $\xi < \eta'$ , whence  $f_{\sigma'}(\lambda) \leq f_\sigma(\lambda)$ . The argument for case 3 is similar to that for case 2. Part b is immediate from part a. For part c let  $\tau_1 = \sigma_{\leq \tau}$  and use part b. Part d follows by parts a and c. □

A scheme  $\sigma \downarrow \lambda$  over  $\lambda$ , for  $\lambda \notin T_\sigma$ , may be defined recursively as follows. Write  $\sigma \downarrow \lambda$  as  $\langle \sigma', \phi' \rangle$ .

- 0: The scheme with  $\sigma' = 0$ .
- 1:  $\tau \downarrow \lambda$  with  $\tau'$  replaced by  $\tau' + 1$ .
- 2:  $\sqcup_{\xi < \eta} \sigma_\xi \downarrow \lambda$ , with  $\phi'(\sigma')(\xi)$  set to  $(\phi(\sigma)(\xi))'$ .
- 3:  $\sqcup_{\xi < \lambda} \sigma_\xi \downarrow \lambda$ , with  $\phi'(\sigma')(\xi)$  set to  $(\phi(\sigma)(\xi))'$ .

**Lemma 2.** *Suppose  $\sigma$  is a scheme,  $\lambda \in \text{In}_\kappa$ , and  $\lambda \notin T_\sigma$ .*

- a,  $f_\sigma \upharpoonright \lambda = f_{\sigma \downarrow \lambda}$ .

b,  $f_\sigma(\lambda) = \sigma \downarrow \lambda$ .

*Proof.* This is lemma 5 of [5]; for convenience the details are given. For part a, the cases of the induction are as follows, where  $\mu \in \text{In}_\lambda$ .

- 0:  $f_0(\mu) = 0 = f_{0 \downarrow \lambda}(\mu)$ .
- 1:  $f_\sigma(\mu) = f_\tau(\mu) + 1 = f_{\tau \downarrow \lambda}(\mu) + 1 = f_{\sigma \downarrow \lambda}(\mu)$ .
- 2:  $f_\sigma(\mu) = \sup_{\xi < \eta} f_{\sigma_\xi}(\mu) = \sup_{\xi < \eta} f_{\sigma_\xi \downarrow \lambda}(\mu) = f_{\sigma \downarrow \lambda}(\mu)$ .
- 3:  $f_\sigma(\mu) = \text{dsup}_{\xi < \kappa} f_{\sigma_\xi}(\mu) = \text{dsup}_{\xi < \lambda} f_{\sigma_\xi \downarrow \lambda}(\mu) = f_{\sigma \downarrow \lambda}(\mu)$ .

For part b, the cases of the induction are as follows.

- 0:  $f_0(\lambda) = 0 = 0 \downarrow \lambda$ .
- 1:  $f_\sigma(\lambda) = f_\tau(\lambda) + 1 = \tau \downarrow \lambda + 1 = \sigma \downarrow \lambda$ .
- 2:  $f_\sigma(\lambda) = \sup_{\xi < \eta} f_{\sigma_\xi}(\lambda) = \sup_{\xi < \eta} \sigma_\xi \downarrow \lambda = \sigma \downarrow \lambda$ .
- 3:  $f_\sigma(\lambda) = \text{dsup}_{\xi < \kappa} f_{\sigma_\xi}(\lambda) = \text{dsup}_{\xi < \lambda} \sigma_\xi \downarrow \lambda = \sigma \downarrow \lambda$ .

□

## 2. Infinitary Veblen Functions

The infinitary Veblen functions were defined in [11]. Further work on them includes [9] and [10]. A certain set of symbols for them are called Klammer-symbols, or Schutte brackets. For convenience a self-contained treatment will be given here, following [11].

For basic facts about normal functions  $f : \text{Ord} \mapsto \text{Ord}$ , and club subclasses  $R \subseteq \text{Ord}$ , see [8], section 4.2. In particular, the following hold.

- If  $f$  is a normal function then  $\text{Ran}(f)$  is a club subclass.
- If  $R$  is a club subclass then the function  $\text{Enum}(R)$  enumerating  $R$  in natural order is normal.
- If  $f$  is normal then the function  $\text{Fix}(f)$  enumerating the subclass  $\{\alpha : f(\alpha) = \alpha\}$  of fixed points of  $f$  is normal.
- The intersection of a set of club class is club.
- Given a normal function  $f$ , suppose  $f(\alpha) > \alpha$ . Let  $\alpha_0 = \alpha$ , and for  $n \geq 0$  let  $\alpha_{n+1} = f(\alpha_n)$ . Then  $\sup_n \alpha_n$  is the smallest fixed point of  $f$  which is greater than  $\alpha$ .
- If  $f(0) > 0$  then  $\text{Fix}(f)(0) > f(0)$ .

Let  $\mathcal{A}$  be the set of ordinal valued sequences  $\alpha_0, \dots, \alpha_\mu$  for some  $\mu$ , where only finitely many  $\alpha_\xi$  are nonzero, and  $\alpha_\mu > 0$  if  $\mu > 0$ .  $\mathcal{A}$  may be ordered in reverse lexicographic order, where if  $\bar{\alpha} = \alpha_0, \dots, \alpha_\mu$  and  $\bar{\alpha}' = \alpha'_0, \dots, \alpha'_{\mu'}$ , then  $\bar{\alpha}' <_{\text{rl}} \bar{\alpha}$  iff  $\mu' < \mu$ , or  $\mu' = \mu$ , and for some  $\gamma \leq \mu$ ,  $\alpha'_\gamma < \alpha_\gamma$  and  $\alpha'_\xi = \alpha_\xi$  for  $\gamma < \xi \leq \mu$ . This relation is well-known to be a well-order (see [11] for example).

Let  $\mathcal{A}_*$  be the sequences obtained from the sequences of  $\mathcal{A}$ , by replacing some  $\alpha_\gamma$  by  $*$ , provided  $\alpha_\delta = 0$  for  $\delta < \gamma$ . The order  $<_{r1}$  may be extended to  $\mathcal{A} \cup \mathcal{A}_*$ , by taking  $\alpha < *$  for  $\alpha \in \text{Ord}$ .

Let  $\theta$  be a cardinal. A function  $\phi : \mathcal{A} \mapsto \text{Ord}$  will be defined, together with functions  $\phi_{\tilde{\alpha}} : \text{Ord} \mapsto \text{Ord}$  for each  $\tilde{\alpha} \in \mathcal{A}_*$ . The cases of the recursion are as follows.

1.  $\phi_{\tilde{\alpha}}(\zeta) = \phi(S_1(\tilde{\alpha}, \zeta))$ , where in  $S_1(\tilde{\alpha}, \zeta)$ ,  $*$  is replaced by  $\zeta$  (and trailing 0's removed if  $\tilde{\alpha}_\mu = *$  and  $\zeta = 0$ ).
2. For  $\alpha \in \text{Ord}$ ,  $\phi(\alpha) = \theta^\alpha$ .
3. For  $\tilde{\alpha} \in \mathcal{A}_*$  with  $\mu > 0$ ,  $\tilde{\alpha}_0 = *$ ,  $\tilde{\alpha}_\gamma = 0$  for  $0 < \gamma < \nu$ , and  $\tilde{\alpha}_\nu > 0$ ,  $\phi_{\tilde{\alpha}} = \text{Enum}(\cap_{\beta < \alpha_\nu, \gamma < \nu} \text{Ran}(\text{Fix}(\phi_{S_2(\alpha, \beta, \gamma)})))$ , where in  $S_2(\alpha, \beta, \gamma)$ ,  $\alpha_0$  is replaced by 0 and  $\alpha_\gamma$  by  $*$  (this is a null operation if  $\gamma = 0$ ), and  $\alpha_\nu$  is replaced by  $\beta$  (and trailing 0's removed if  $\nu = \mu$  and  $\beta = 0$ ).

**Lemma 3.** For  $\tilde{\alpha} \in \mathcal{A}_*$ ,  $\phi_{\tilde{\alpha}}$  is normal.

*Proof.* The proof is by induction on  $<_{r1}$ . For the basis,  $\tilde{\alpha} = *$ ; the lemma follows since  $\zeta \mapsto \theta^\zeta$  is normal. For the induction step, if  $\tilde{\alpha}_0 = *$  the lemma follows by the induction hypothesis and the definition of  $\phi$ . Otherwise, using obvious notation write  $\tilde{\alpha}$  as  $0\bar{0}*\bar{\eta}$ ; then  $\phi_{0\bar{0}*\bar{\eta}}(\zeta) = \phi_{*\bar{0}\zeta\bar{\eta}}(0)$ . Suppose  $\xi < \zeta$ . Directly from the definition of  $\phi$ ,  $\text{Ran}(\phi_{*\bar{0}\xi\bar{\eta}}) \supseteq \text{Ran}(\phi_{*\bar{0}\zeta\bar{\eta}})$  and  $\phi_{*\bar{0}\xi\bar{\eta}}(0) < \phi_{*\bar{0}\zeta\bar{\eta}}(0)$ . In particular,  $\phi_{0\bar{0}*\bar{\eta}}(\xi) < \phi_{0\bar{0}*\bar{\eta}}(\zeta)$ ; that is,  $\phi_{\tilde{\alpha}}$  is increasing.

For  $\zeta \in \text{Lim}$  let  $\chi = \sup_{\xi < \zeta} \phi_{\tilde{\alpha}}(\xi)$ . By what was just proved,  $\chi \leq \phi_{\tilde{\alpha}}(\zeta)$ , whence  $\chi \leq \phi_{*\bar{0}\zeta\bar{\eta}}(0)$ . Thus, to show that  $\phi_{\tilde{\alpha}}$  is continuous it suffices to show that  $\chi \in \text{Ran}(\phi_{*\bar{0}\zeta\bar{\eta}})$ , since then  $\chi = \phi_{*\bar{0}\zeta\bar{\eta}}(0) = \phi_{\tilde{\alpha}}(\zeta)$ . For this in turn it suffices to show that  $\phi_{S_2(*\bar{0}\zeta\bar{\eta}, \beta, \gamma)}(\chi) = \chi$  for  $\beta < \zeta$ ,  $\gamma < \nu$ .

Suppose  $\beta < \xi$ , and for ease of notation let

$$\tilde{\alpha}_1 \text{ denote } S_2(*\bar{0}\zeta\bar{\eta}, \beta, \gamma), \text{ and } \tilde{\alpha}_2 \text{ denote } *\bar{0}\zeta\bar{\eta}.$$

From the definition of  $\phi$ ,  $\text{Ran}(\phi_{\tilde{\alpha}_2}) \subseteq \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}_1}))$ , and so  $\phi_{\tilde{\alpha}_2}(0) \in \text{Fix}(\phi_{\tilde{\alpha}_1})$ , whence  $\phi_{\tilde{\alpha}_1}(\phi_{\tilde{\alpha}_2}(0)) = \phi_{\tilde{\alpha}_2}(0)$ . Using this fact,

$$\begin{aligned} \phi_{S_2(\phi_{*\bar{0}\zeta\bar{\eta}}, \beta, \gamma)}(\chi) &= \\ \phi_{S_2(\phi_{*\bar{0}\zeta\bar{\eta}}, \beta, \gamma)}(\sup_{\xi < \zeta} \phi_{\tilde{\alpha}}(\xi)) &= \sup_{\xi < \zeta} \phi_{S_2(\phi_{*\bar{0}\zeta\bar{\eta}}, \beta, \gamma)}(\phi_{\tilde{\alpha}}(\xi)) = \\ \sup_{\beta < \xi < \zeta} \phi_{S_2(\phi_{*\bar{0}\zeta\bar{\eta}}, \beta, \gamma)}(\phi_{\tilde{\alpha}}(\xi)) &= \sup_{\beta < \xi < \zeta} \phi_{\tilde{\alpha}}(\xi) = \chi. \end{aligned} \quad \square$$

**Lemma 4.** If  $\tilde{\beta} <_{r1} \tilde{\alpha}$  then  $\text{Ran}(\text{Fix}(\phi_{\tilde{\beta}})) \cup \{1\} \supseteq \text{Ran}(\phi_{\tilde{\alpha}})$ .

*Proof.* Writing  $\tilde{\alpha}$  as  $0\bar{0}*\bar{\alpha}$ ,  $\phi_{0\bar{0}*\bar{\alpha}}(\zeta) = \phi_{*\bar{0}\zeta\bar{\alpha}}(0)$ , provided  $\zeta \neq 0$  if  $\bar{\alpha}$  is null. By the definition of  $\phi$  and induction on  $<_{r1}$ ,  $\phi_{\tilde{\alpha}}(\zeta) \in \text{Ran}(\text{Fix}(\phi_{\tilde{\beta}})) \cup \{1\}$ , proving the lemma. □

Let  $\mathcal{A}_*^0 = \{\tilde{\alpha} \in \mathcal{A}_* : \tilde{\alpha}_0 = *\}$ . The following are readily seen, for  $\tilde{\beta} <_{r1} \tilde{\alpha}$ .

- For  $\tilde{\alpha} \in \mathcal{A}_*^0$ ,  $1 \in \text{Ran}(\phi_{\tilde{\alpha}})$  iff  $\tilde{\alpha} = *$ .
- For  $\tilde{\alpha}, \tilde{\beta} \in \mathcal{A}_*^0$ , If  $\tilde{\beta} <_{\text{rl}} \tilde{\alpha}$  then  $\text{Ran}(\text{Fix}(\phi_{\tilde{\beta}})) \supseteq \text{Ran}(\phi_{\tilde{\alpha}})$ .
- For  $\tilde{\alpha}, \tilde{\beta} \in \mathcal{A}_*$ ,  $\text{Ran}(\text{Fix}(\phi_{\tilde{\beta}})) \supseteq \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}}))$ .

For  $\mu \in \text{Ord}$  let  $\tilde{\alpha}_\mu$  be the element of  $\mathcal{A}_*$  where  $\tilde{\alpha}_{\mu 0} = *$ ,  $\tilde{\alpha}_{\mu\mu} = 1$ , and  $\tilde{\alpha}_{\mu\xi} = 0$  for  $0 < \xi < \mu$ . Let  $\psi$  be the function  $\mu \mapsto \tilde{\alpha}_\mu(0)$ .

**Lemma 5.** *The function  $\psi$  is normal.*

*Proof.* For  $\mu \in \text{Ord}$  let  $\tilde{\alpha}_\mu^*$  be the element of  $\mathcal{A}_*$  where  $\tilde{\alpha}_{\mu\mu} = *$  and  $\tilde{\alpha}_{\mu\xi} = 0$  for  $0 \leq \xi < \mu$ . From the definition of  $\phi$ ,  $\phi_{\tilde{\alpha}_\mu^*} = \text{Enum}(\cap_{\gamma < \mu} \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}_\gamma^*}))$ ). It follows from this and lemma 4 that  $\mu \mapsto \tilde{\alpha}_\mu(0)$  is increasing. For  $\mu \in \text{Lim}$  let  $\chi = \sup_{\gamma < \mu} \phi_{\tilde{\alpha}_\gamma}(0)$ . From the definition of  $\phi$ , for  $\gamma < \delta$   $\text{Ran}(\phi_{\tilde{\alpha}_\delta}) \subseteq \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}_\gamma^*}))$ . Using this fact, it follows by arguments as given above that  $\phi_{\tilde{\alpha}_\mu}(0) = \chi$ .  $\square$

The least fixed point of  $\psi$  is called the “large Veblen ordinal”. Notation varies;  $\Lambda_0$  will be used here. Clearly,  $\Lambda_0 = \phi_{\tilde{\alpha}_{\Lambda_0}}(0) = \sup_{\mu < \Lambda_0} \phi_{\tilde{\alpha}_\mu}(0)$ . It follows that for  $\mu < \Lambda_0$ ,  $\Lambda_0$  is a fixed point of  $\phi_{\tilde{\alpha}_\mu}$ . Using lemma 4, if  $\tilde{\alpha} < \tilde{\alpha}_{\Lambda_0}$  then  $\Lambda_0$  is a fixed point of  $\phi_{\tilde{\alpha}}$ , whence if  $\zeta < \Lambda_0$  then  $\phi_{\tilde{\alpha}}(\zeta) < \Lambda_0$ .

Say that  $\bar{\alpha} \in \mathcal{A}$  is maximal if  $\phi(\tilde{\beta}) \neq \phi(\bar{\alpha})$  whenever  $\beta >_{\text{rl}} \alpha$ .

**Lemma 6.** *For  $\bar{\alpha} \in \mathcal{A}$ , let  $\nu_1 > \dots > \nu_k$  be the indices such that  $\alpha_{\nu_i} > 0$ .*

- a. *For  $i > 1$ ,  $\phi(\bar{\alpha}) > \alpha_{\nu_i}$ .*
- b.  *$\phi(\bar{\alpha}) > \alpha_{\nu_1}$  iff  $\bar{\alpha}$  is maximal.*
- c. *If  $\phi(\bar{\alpha}) \notin \text{Ran}(\psi)$  then  $\nu_k < \phi(\bar{\alpha})$ .*
- d. *If  $\tau \in \text{Ran}(\phi_*)$  then  $\tau = \phi(\bar{\alpha})$  for a unique maximal  $\bar{\alpha}$ .*

*Proof.* For  $1 \leq j < i$  let  $\bar{\alpha}_i$  be  $\bar{\alpha}$  with  $\alpha_{\nu_j}$  replaced by 0. It is easily seen that  $\phi(\bar{\alpha}_{i-1}) > \phi(\bar{\alpha}_i) \geq \alpha_{\nu_i}$  for  $i > 1$ , and part a follows. Write  $\bar{\alpha}$  as  $\bar{0}\alpha_\nu\alpha_{\nu+1}\bar{\alpha}'$ , where for ease of notation  $\mu > \nu$  is assumed. Let  $\tilde{\alpha} = \bar{0}*\alpha_{\nu+1}\bar{\alpha}'$ . Since  $\phi_{\tilde{\alpha}}$  is normal,  $\phi(\alpha) \geq \alpha_\nu$ . Let  $\tilde{\beta} = *\bar{0}0, \alpha_{\nu+1} + 1, \bar{\alpha}'$ . Then  $\text{Ran}(\phi_{\tilde{\beta}}) = \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}}))$ , and so  $\phi(\bar{\alpha}) = \alpha_\nu$  iff  $\alpha_\nu \in \text{Ran}(\phi_{\tilde{\beta}})$  iff  $\bar{\alpha}$  is not maximal. The case  $\nu = \mu$  is similar, proving part b. For part c, let  $\gamma \in \text{Ran}(\psi)$  be least such that  $\phi(\bar{\alpha}) < \gamma$ . Then  $\phi(\bar{\alpha}_{\nu_k}) \leq \phi(\bar{\alpha}) < \gamma$ , whence  $\nu_k < \phi(\bar{\alpha}_{\nu_k}) \leq \phi(\bar{\alpha})$ . For part d, there is a  $\mu$  such that  $\tau < \phi_{\tilde{\alpha}_\mu}(0) \leq \phi_{\tilde{\alpha}_\mu}(\tau)$ , and so there is a lexicographically least  $\tilde{\alpha}$  such that  $\tau < \phi_{\tilde{\alpha}}(\tau)$ .

If  $\tilde{\alpha} = *$  then  $\tau = \phi(\zeta)$  for some  $\zeta$  by hypothesis, and  $\zeta < \tau$  since  $\zeta = \tau$  is impossible. Now,  $\tau = \phi_{\tilde{\beta}}(\tau)$  for  $\tilde{\beta} <_{\text{rl}} \tilde{\alpha}$ . If  $\tilde{\alpha}_0 = *$  then by the definition of  $\phi$   $\tau = \phi_{\tilde{\alpha}}(\zeta)$  for some  $\zeta$ , and  $\zeta = \tau$  is impossible. Otherwise  $\tilde{\alpha} = 0\bar{0} * \bar{\alpha}'$ , and  $\tau < \phi_{0\bar{0}*\bar{\alpha}'}(\tau) = \phi_{*\bar{0}\tau\bar{\alpha}'}(0)$ ; but then by the definition of  $\phi$   $\tau = \phi_{*\bar{0}\tau\bar{\alpha}'}(\zeta)$  for some  $\zeta$ , which is impossible.  $\square$

An infinitary Veblen function is a function  $\phi_k(\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k)$  where  $k \geq 1$ . If  $\xi_1 < \dots < \xi_k$  and  $\alpha_i > 0$  the arguments are said to be proper; the value of  $\phi_k$  is  $\phi(\bar{\alpha})$  where  $\bar{\alpha}$  is the sequence  $\zeta, \dots, \alpha_1 \dots \alpha_2 \dots \alpha_k$ , where  $\alpha_i$  occurs in position  $1 + \zeta_i$ , and all other elements of  $\bar{\alpha}$  are 0. If the arguments are improper the value is 0.

**Corollary 7.** *For  $\tau \in \text{Ran}(\phi_*)$  with  $\tau < \Lambda_0$  there are a unique  $k$  and values  $\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k < \tau$  such that  $\tau = \phi_k(\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k)$ .*

*Proof.* Immediate. □

Let  $O$  denote the function with ordinal arguments, where  $O(\alpha_1, \alpha_2)$  equals  $-1$  if  $\alpha_1 < \alpha_2$ ,  $0$  if  $\alpha_1 = \alpha_2$ , and  $+1$  if  $\alpha_1 > \alpha_2$ ,

**Lemma 8.** *Suppose  $\tilde{\alpha}_1, \tilde{\alpha}_2 \in \mathcal{A}_*^0$  and  $\tilde{\alpha}_1 <_{rl} \tilde{\alpha}_2$ . Then  $O(\phi_{\tilde{\alpha}_1}(\zeta_1), \phi_{\tilde{\alpha}_2}(\zeta_2)) = O(\zeta_1, \phi_{\tilde{\alpha}_2}(\zeta_2))$ .*

*Proof.*  $\text{Ran}(\phi_{\tilde{\alpha}_2}) \subseteq \text{Ran}(\text{Fix}(\phi_{\tilde{\alpha}_1}))$ , so  $\phi_{\tilde{\alpha}_1}(\phi_{\tilde{\alpha}_2}(\zeta_2)) = \phi_{\tilde{\alpha}_2}(\zeta_2)$ ; the lemma follows. □

### 3. $IV_o$ Terms

As in [7], a CNF function is a function  $C_k(\eta_k, \sigma_k, \dots, \eta_1, \sigma_1)$ , where  $k \geq 1$ . If  $\eta_k > \dots > \eta_1$  and  $0 < \sigma_i < \theta$  the arguments are said to be proper, and the value of  $C_k$  is  $\theta^{\eta_k} \cdot \sigma_k + \dots + \theta^{\eta_1} \cdot \sigma_1$ . If the arguments are improper, the value is defined to be 0.

Let  $IV_o$  be the system of terms, whose leaves are ordinals  $\sigma$ , and whose interior nodes are either a CNF function  $C_k$  or an infinitary Veblen function  $\phi_k$ . A proper term is one where the arguments are proper at each interior node. A normal form term is a proper term where at each interior node, the value of the function is greater than the value of any of its arguments.

**Theorem 9.** *If  $\alpha < \Lambda_0$  then there is a unique normal form  $IV_o$  term whose value is  $\alpha$ .*

*Proof.* Let  $\kappa^{+\eta_k} \cdot \sigma_k + \dots + \kappa^{+\eta_1} \cdot \sigma_1$  be the Cantor normal form for  $\alpha$ . The proof is by induction on  $\alpha$ , with cases as follows, where  $t$  denotes the term.

Case 0:  $\alpha = 0$ . Then  $t = 0$ .

Case 1:  $k = 1$  and  $\eta_1 = 0$ . Then  $t = \sigma_1$ .

Case 2:  $k > 1$  or  $\sigma_1 > 1$  or  $\eta_1 > 0$  and  $\eta_1 < \kappa^{+\eta}$ . Then  $t = C_k(t_{\eta_k}, t_{\sigma_k}, \dots, t_{\eta_1}, t_{\sigma_1})$ .

Case 3:  $k = 1$  and  $\sigma_1 = 1$  and  $\eta_1 = \kappa^{+\eta}$ . In this case,  $t = \phi_k(t_\zeta, t_{\xi_1}, t_{\alpha_1}, \dots, t_{\xi_k}, t_{\alpha_k})$  where  $\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k < \alpha$  are the unique values as in lemma 7.  $\square$

If  $t_i$  is a term with value  $\alpha_i$  for  $i = 1, 2$ , write  $O(t_1, t_2)$  for  $O(\alpha_1, \alpha_2)$ . Let  $P(t)$  equal 1 if  $t$  is proper, else 0. Let  $N(t)$  equal 1 if  $t$  is normal, else 0.

- Lemma 10.**
- a.  $O(t_1, t_2)$  depends only on the values  $O(\sigma_1, \sigma_2)$  and  $P(\sigma)$  for the leaves of  $t_1, t_2$ .
  - b.  $P(t)$  depends only on the values  $O(\sigma_1, \sigma_2)$  and  $P(\sigma)$  for the leaves of  $t$ .
  - c.  $N(t)$  depends only on the values  $O(\sigma_1, \sigma_2)$  and  $P(\sigma)$  for the leaves of  $t$ .

*Proof.* Parts a and b are proved by induction on the total number of nodes. For part b, if  $t$  is a leaf then  $P(t) = 1$ . If  $t$  is a CNF function,  $P(t)$  iff  $P(s)$  holds for all subterms  $s$ , and certain inequalities hold among the subterms. If  $t$  is a IV function the argument is similar.

Part a falls into the following cases. If  $t_1$  a leaf and  $t_2$  a leaf then  $O(t_1, t_2)$  is known by hypothesis. Suppose  $t_1$  an ordinal and  $t_2$  is a CNF function. If  $t_2$  is improper then  $O(t_1, t_2)$  equals 0 if  $t_1 = 0$ , else 1. If  $t_2$  is proper then  $O(t_1, t_2)$  equals  $-1$ , unless  $t_2 = \kappa^{+0} \cdot \sigma_2$ , in which case it equals  $O(t_1, \sigma_2)$ . Suppose  $t_1$  an ordinal and  $t_2$  is an IV function. If  $t_2$  is improper then  $O(t_1, t_2)$  equals 0 if  $t_1 = 0$ , else 1. If  $t_2$  is proper then  $O(t_1, t_2) = -1$ . In the remaining cases, if  $t_1$  and  $t_2$  are both improper then  $O(t_1, t_2) = 0$ , if  $t_1$  is improper and  $t_2$  is proper then  $O(t_1, t_2) = -1$ , and if  $t_1$  is proper and  $t_2$  is improper then  $O(t_1, t_2) = 1$ ; so it may be assumed that  $t_1$  and  $t_2$  are proper. If  $t_1$  and  $t_2$  are both CNF functions,  $O(t_1, t_2)$  is determined by the lexicographic order, using the induction hypothesis. Suppose  $t_1$  is a CNF function and  $t_2$  is an IV functions. If  $\eta_k < \alpha_2$  (where  $\eta_k$  is the leading exponent in  $t_1$ ) then  $\alpha_1 < \alpha_2$ . If  $\eta_k > \alpha_2$  then  $\alpha_1 > \alpha_2$ . If  $\eta_k = \alpha_2$  then  $\alpha_1 > \alpha_2$ , unless  $k = 1$  and  $\sigma_k = 1$ , in which case  $\alpha_1 = \alpha_2$ . If  $t_1$  is an IV function and  $t_2$  is a CNF functions, reverse the roles of  $t_1$  and  $t_2$  in the preceding argument. Suppose  $t_1$  and  $t_2$  are IV functions, say  $\phi_{\tilde{\alpha}_1}(\zeta_1)$  and  $\phi_{\tilde{\alpha}_2}(\zeta_2)$ , If  $\tilde{\alpha}_1 <_{r1} \tilde{\alpha}_2$  (which may be determined from the arguments using the induction hypothesis) then  $O(t_1, t_2) = O(\zeta_1, \alpha_2)$ . If  $\tilde{\alpha}_1 = \tilde{\alpha}_2$  then  $O(t_1, t_2) = O(\zeta_1, \zeta_2)$ . If  $\tilde{\alpha}_2 <_{r1} \tilde{\alpha}_1$  then  $O(t_1, t_2) = O(\alpha_1, \zeta_2)$ .

Part c now follows inductively. The claim is trivial if  $t$  is a leaf. If  $t$  is a CNF function then  $t$  is normal if its subterms are normal, except if  $k = 1$  and  $\eta_k = 0$ . If  $t$  is an IV function then  $t$  is normal if its subterms are normal, except if  $t = \phi_{\tilde{\alpha}_1}(\phi_{\tilde{\alpha}_2}(\zeta_2))$  where  $\tilde{\alpha}_1 <_{r1} \tilde{\alpha}_2$ .  $\square$

### 4. IV Scheme Terms

From hereon the cardinal  $\theta$  will be  $\kappa^+$  for some  $\kappa \in \text{Inac}$ . A scheme term over  $\kappa$  is defined to be a term, whose leaves are 0 or a scheme over  $\kappa$  with  $\sigma > 0$ ; and whose interior nodes are either a CNF function  $C_k$  or an infinitary Veblen function  $\phi_k$  where  $k \geq 1$ . The value  $\alpha$  of a term  $\alpha$  is the value of the  $\text{IV}_o$  term obtained by replacing each scheme  $\sigma$  by its value  $\sigma$ . Clearly  $\alpha$  depends only on the values of the schemes.

It follows from theorem 9 that for each  $\alpha < \Lambda_0$  there is a normal form term  $\alpha$ , whose ordinal is  $\alpha$ .  $\alpha$  is unique, up to the choice of schemes  $\sigma$  for the leaf ordinals  $\sigma$ . From hereon, unless specifically stated otherwise, IV scheme terms will be assumed to be in normal form.

For a term  $\alpha$  let  $T_\alpha$  be the union over the  $\sigma$  occurring at leaves of  $\alpha$  of the  $T_\sigma$ .

Let  $\alpha \downarrow \lambda$  be  $\alpha$ , with each leaf  $\sigma$  replaced by  $\sigma \downarrow \lambda$ , each node  $C_k$  replaced by  $C_{\lambda k}$ , and each node  $\phi_k$  replaced by  $\phi_{\lambda k}$ . Given  $\alpha$ , the notation  $\alpha \downarrow \lambda$  will be used to denote the ordinal specified by  $\alpha \downarrow \lambda$ .

**Lemma 11.** *Suppose  $\alpha$  is a scheme,  $\lambda \in \text{In}_\kappa$ , and  $\lambda \notin T_\alpha$ .*

- a,  $f_\alpha \upharpoonright \lambda = f_{\alpha \downarrow \lambda}$ .
- b,  $f_\alpha(\lambda) = \alpha \downarrow \lambda$ .

*Proof.* For schemes this is lemma 2. Let the argument list of a function be denoted  $A_1, \dots, A_l$ . For part a, suppose  $\mu < \lambda$ . For CNF functions,  $f_{\alpha \downarrow \lambda}(\mu) = f_{C_k(A_1 \downarrow \lambda, \dots, A_l \downarrow \lambda)}(\mu) = C_{\mu k}(f_{A_1 \downarrow \lambda}(\mu), \dots, f_{A_l \downarrow \lambda}(\mu)) = C_{\mu k}(f_{A_1}(\mu), \dots, f_{A_l}(\mu)) = f_{C_k(A_1, \dots, A_l)}(\mu) = f_\alpha(\mu)$ . The proof for IV functions is similar. For part b, for CNF functions,  $f_\alpha(\lambda) = f_{C_k(A_1, \dots, A_l)}(\lambda) = C_{\lambda k}(f_{A_1}(\lambda), \dots, f_{A_l}(\lambda)) = C_{\lambda k}(A_1 \downarrow \lambda, \dots, A_l \downarrow \lambda) = C_k(A_1, \dots, A_l) \downarrow \lambda = \alpha \downarrow \lambda$ . The proof for IV functions is similar. □

### 5. Function Chains for IV Terms

For each IV term  $\alpha$  a function  $f_\alpha : \text{In}_\kappa \mapsto \kappa$  may be defined by recursion on  $\alpha$ . The cases of the recursion are numbered as in theorem 9, and are as follows. In cases 2 and 3,  $C_{\lambda k}$  and  $\phi_{\lambda k}$  is written for the function over the cardinal  $\lambda^+$ .

- 0.  $f_0 = 0$
- 1.  $f_\sigma$  where  $\alpha$  is the scheme  $\sigma$ .
- 2.  $f_\alpha(\lambda) = C_{\lambda k}(f_{\eta_k}(\lambda), f_{\sigma_k}(\lambda), \dots, f_{\eta_1}(\lambda), f_{\sigma_1}(\lambda))$ .
- 3.  $f_\alpha(\lambda) = \phi_{\lambda k}(f_\zeta(\lambda), f_{\xi_1}(\lambda), f_{\alpha_1}(\lambda), \dots, f_{\xi_k}(\lambda), f_{\alpha_k}(\lambda))$ .



**Theorem 12.** Suppose  $\alpha$ , etc. are proper terms, and  $\lambda \in \text{In}_\kappa$ .

- a. If  $\lambda \notin T_{\alpha_1} \cup T_{\alpha_2}$  then  $O(f_{\alpha_1}(\lambda), f_{\alpha_2}(\lambda)) = O(\alpha_1, \alpha_2)$ .
- b. If  $\lambda \notin T_\alpha$  then  $P(f_\alpha(\lambda)) = P(\alpha)$ .
- c. If  $\lambda \notin T_\alpha$  then  $N(f_\alpha(\lambda)) = N(\alpha)$ .

*Proof.* This follows by theorem 10 and lemma 1. □

**Theorem 13.** Suppose  $\alpha$  and  $\beta$  are IV terms,  $\alpha = \beta + 1$ ,  $\lambda \in \text{In}_\kappa$ , and  $\lambda \notin T_\alpha \cup T_\beta$ . Then  $f_\alpha(\lambda) = f_\beta(\lambda) + 1$ .

*Proof.* In the CNF for  $\alpha$ ,  $\eta_1 = 0$ ; let  $\alpha_p$  be  $\alpha$  with the last term  $\sigma_1$  replaced by  $\tau_1$ , or removed entirely if  $\sigma_1 = 1$ . By theorem 12,  $f_\beta(\lambda) = Tf_{\alpha_p}(\lambda)$ . It follows that  $f_\beta(\lambda) + 1 = Tf_{\alpha_p}(\lambda) + 1 = f_\alpha(\lambda)$ . □

### 6. Set Chains for IV Terms

The reader is assumed to be familiar with [7]. A definition of  $H_\alpha$  for an IV term  $\alpha$  will be given. A definition of  $H_{1\alpha}$  was given in [7] only for a subset of the BV terms.  $H$  is a map from subsets of  $\text{In}_\kappa$  to subsets of  $\text{In}_\kappa$ ; when it is necessary to specify  $\kappa$ ,  $H_\alpha$  will be denoted  $H_{\kappa\alpha}$ .

For an IV term  $\alpha$  and  $X \subseteq \text{In}_\kappa$ , say that  $\lambda \in H_\alpha(X)$  iff  $\lambda \in X$  and  $H_\beta(X \cap \lambda)$  is a stationary subset of  $\lambda$  for all IV terms  $\beta$  over  $\lambda$  where  $\beta < f_\alpha(\lambda)$ .

For  $X, Y \subseteq \text{In}_\kappa$  say that  $X \subseteq_t Y$  if  $\{\lambda \in \text{In}_\kappa : \lambda \in X \text{ and } \lambda \notin Y\}$  is thin; and similarly for  $\equiv_t$ .

**Theorem 14.** If  $\alpha' \leq \alpha$  then for any  $X \subseteq \text{In}_\kappa$ ,  $H_{\alpha'}(X) \supseteq_t H_\alpha(X)$ .

*Proof.* The proof is by induction on  $\kappa$ . For the basis,  $\kappa$  is the smallest inaccessible cardinal,  $X$  is always empty, and the claim is trivial. For arbitrary  $\kappa$ , there is a thin set  $T$  such that if  $\lambda \in \text{In}_\kappa$  and  $\lambda \notin T$  then  $f_{\alpha'}(\lambda) \leq f_\alpha(\lambda)$ . For such  $\lambda$ , if  $H_\beta(X \cap \lambda)$  is stationary for  $\beta < f_\alpha(\lambda)$ , then by the induction hypothesis  $H_\beta(X \cap \lambda)$  is stationary for  $\beta < f_{\alpha'}(\lambda)$ . □

**Lemma 15.** Suppose  $\alpha$  is an IV term,  $X \subseteq \text{In}_\kappa$ ,  $\lambda \in \text{In}_\kappa$ , and  $\lambda \notin T_\alpha$ . Then  $H_{\alpha \downarrow \lambda}(X \cap \lambda) = H_\alpha(X) \cap \lambda$ .

*Proof.* Suppose  $\mu < \lambda$ . Then  $\mu \in H_{\alpha \downarrow \lambda}(X \cap \lambda)$  iff  $H_\beta(X \cap \mu)$  is stationary for  $\beta < f_{\alpha \downarrow \lambda}(\mu)$  iff (by lemma 11)  $H_\beta(X \cap \mu)$  is stationary for  $\beta < f_\alpha(\mu)$  iff  $\mu \in H_\alpha(X)$ . □

**Theorem 16.** *Suppose  $\alpha = \beta + 1$ ,  $\lambda \in \text{In}_\kappa$ , and  $\lambda \notin T_\alpha$ . Then  $\lambda \in H_\alpha(X)$  iff  $\lambda \in H(H_\beta(X))$ .*

*Proof.*  $\lambda \in H_\alpha(X)$  iff  $\lambda \in X$  and  $H_\gamma(X \cap \lambda)$  is stationary for  $\gamma$  with  $\gamma < f_\alpha(\lambda)$  iff  $\lambda \in H_\beta(X)$  and (A)  $H_\gamma(X \cap \lambda)$  is stationary for  $\gamma$  with  $\gamma = f_\alpha(\lambda)$ . Using lemma 11.b, (A) holds iff  $H_{\beta \downarrow \lambda}(X \cap \lambda)$  is stationary, and by lemma 15 this holds iff  $H_\beta(X) \cap \lambda$  is stationary.  $\square$

Recall from [7] the definition of  $\rho_R$ . By theorems 14 and 16, it follows that if  $H_\alpha(\text{In}_\kappa)$  is stationary then  $\rho_R(H_\alpha(\text{In}_\kappa)) \geq \alpha$ . This was proved in [7] only for a subset of the BV terms (with a different definition of  $H_\alpha$ ; but see below).

### 7. A Continuity Property

As noted in [11], explicit ascending sequences can be given for limit ordinals less than  $\Lambda_0$ , using IV terms. A sequences  $\alpha_\xi$  for  $\alpha$  may be defined by recursion on  $\alpha$ . Case 0 is irrelevant. In case 1, let  $\sigma_\xi$  be as determined by the last node of  $\sigma$ .

Case 2 is divided into subcases as follows. In subcase 2.1, in the CNF for  $\alpha$ ,  $k > 1$ . Write  $\alpha = \alpha_1 + \alpha_2$ . Let  $\alpha_\xi$  equal  $\alpha_1 + \alpha_{2\xi}$ .

In subcase 2.2  $\alpha = \kappa^{+\eta} \cdot \sigma$  where  $\sigma \in \text{Lim}$  and  $\sigma > 1$ . Let  $\alpha_\xi$  equal  $\kappa^{+\eta} \cdot \sigma_\xi$ .

In subcase 2.3  $\alpha = \kappa^{+\eta}$  where  $\eta \in \text{Lim}$  and  $\eta < \kappa^{+\eta}$ . Let  $\alpha_\xi$  equal  $\kappa^{+\eta_\xi}$ .

In subcase 2.4  $\alpha = \kappa^{+\eta}$  where  $\eta = \eta^- + 1$ . Let  $\alpha_\xi$  equal  $\kappa^{+\eta^-} \cdot \xi$ , where  $\xi$  is some particular scheme with ordinal  $\xi$ .

To divide case 3 into subcases, write the IV function as  $\phi_k(\zeta, \pi_1, \alpha_1, \dots)$ . Cases will be denoted XYZ where X is the type of  $\zeta$  (L for limit, 0, or S for successor), Y is the type of  $\alpha_1$  (L or S), and Z is the type of  $\xi_1$  (L, 0, or S); in some cases Y or Z may be \*, denoting any possibility. If  $\alpha$  is a successor ordinal write  $\alpha^-$  for its predecessor. Sequences  $\alpha_\xi$  in each case are as follows.

Case L\*\*:  
 $\alpha_\xi = \phi_k(\zeta_\xi, \pi_1, \alpha_1, \dots)$ .

Case 0L\*:  
 $\alpha_\xi = \phi_k(0, \pi_1, \alpha_{1\xi}, \dots)$ .

Case SL\*:  
 $\gamma = \phi_k(\zeta^-, \pi_1, \alpha_1, \dots) + 1, \alpha_\xi = \phi_k(\gamma, \pi_1, \alpha_{1\xi}, \dots)$ .

Case 0SL:  
 $\alpha_\xi = \phi_k(0, \pi_{1\xi}, \alpha_1^-, \dots)$ .

Case SSL:  
 $\gamma = \phi_k(\zeta^-, \pi_1, \alpha_1, \dots) + 1, \alpha_\xi = \phi_k(\gamma, \pi_1, \alpha_{1\xi}, \dots)$ .

Case 0S0:  
 $\alpha_0 = \phi_k(0, 0, \alpha_1^-, \dots), \alpha_{n+1} = \phi_k(\alpha_n, 0, \alpha_1^-, \dots)$ .

Case SS0:  
 $\alpha_0 = \phi_k(\zeta^-, 0, \alpha_1, \dots) + 1, \alpha_{n+1} = \phi_k(\alpha_n, 0, \alpha_1^-, \dots)$ .

Case OSS:  $\alpha_0 = \phi_{k+1}(0, \pi_1^-, 1, \pi_1, \alpha_1^-, \dots)$ ,  $\alpha_{n+1} = \phi_{k+1}(0, \pi_1^-, \alpha_n, \pi_1, \alpha_1^-, \dots)$ .

Case SSS:  $\alpha_0 = \phi_k(\zeta^-, \pi_1, \alpha_1, \dots)$ ,  $\alpha_{n+1} = \phi_{k+1}(0, \pi_1^-, \alpha_n, \pi_1, \alpha_1^-, \dots)$ .

**Lemma 17.** *Suppose  $\alpha$  is a term,  $\lambda \in In_\kappa$ , and  $\lambda \notin T_\alpha$ . If  $\alpha$  is a successor ordinal then  $f_\alpha(\lambda)$  is a successor ordinal. If  $Cf(\alpha) = \eta$  where  $\eta < \kappa$  (whence  $\eta < \lambda$ ) then  $Cf(f_\alpha(\lambda)) = \eta$ . If  $Cf(\alpha) = \kappa$  then  $Cf(f_\alpha(\lambda)) = \kappa$ . If  $Cf(\alpha) = \kappa^+$  then  $Cf(f_\alpha(\lambda)) = \lambda^+$ .*

*Proof.* This follows by induction on  $\alpha$ , using the ascending sequences given above, and the definition of  $f_\alpha(\lambda)$ . □

The notation  $f \leq_t g$  will be used to denote that  $\{\lambda : f(\lambda) > g(\lambda)\}$  is thin; and similarly for  $=_t$ . It is readily verified that sup and dsup produce least upper bounds in the order  $\leq_t$ . See also lemma 1.2 of [2]

**Theorem 18.** *If  $\eta < \kappa$  and  $\alpha_\xi$  for  $\xi < \eta$  is an ascending sequence with  $\alpha = \sup_{\xi < \eta} \alpha_\xi$  then  $f_\alpha =_t \sup_{\xi < \eta} f_{\alpha_\xi}$ .*

*Proof.* The proof is by induction on  $\alpha$ . By remarks preceding the theorem it suffices to prove the claim for a particular sequence  $\alpha_\xi$ . The sequence used will be that given above. Case 0 is irrelevant. Case 1 follows directly. The remaining cases follow inductively, using properties of sup, with subcase 2.4 being irrelevant. □

The theorem can be improved slightly. The set of  $\lambda$  where the equation might not hold is contained in  $T_\alpha \cup (\eta + 1) \cup \cup_{\xi < \eta} T_{\alpha_\xi}$ .

**Theorem 19.** *If  $\alpha_\xi$  for  $\xi < \kappa$  is an ascending sequence with  $\alpha = \sup_{\xi < \kappa} \alpha_\xi$  then  $f_\alpha =_t \text{dsup}_{\xi < \kappa} f_{\alpha_\xi}$ .*

*Proof.* The proof is by induction on  $\alpha$ . As in the preceding theorem it suffices to prove the claim for the sequence given above. Case 0 is irrelevant. Case 1 follows directly. The remaining cases follow inductively, using the definition of dsup and properties of sup, with subcase 2.4 and the subcases of case 3 not involving an L being irrelevant. □

**Theorem 20.** *a. Suppose  $\eta < \kappa$ ,  $\alpha_\xi$  is ascending, and  $\alpha = \sup_{\xi < \eta} \alpha_\xi$ .*

*Then  $H_\alpha \equiv_t \cap_{\xi < \eta} H_{\alpha_\xi}$ .*

*b. Suppose  $\alpha_\xi$  is ascending and  $\alpha = \sup_{\xi < \kappa} \alpha_\xi$ . Then  $H_\alpha \equiv_t$*

*$Bt_{\xi < \kappa} H_{\alpha_\xi}$ .*

*Proof.* For part a,  $\lambda \in H_\alpha(X)$  iff  $H_\beta(X \cap \lambda)$  is stationary whenever  $\beta < f_\alpha(\lambda)$ . By theorem 18, except for a thin set of  $\lambda$ , this holds iff, for all  $\xi < \eta$ ,  $H_\beta(X \cap \lambda)$  is stationary whenever  $\beta < f_{\alpha_\xi}(\lambda)$ , iff  $\lambda \in H_{\alpha_\xi}(X)$ . Part b follows similarly, with “ $\xi < \eta$ ” replaced by “ $\xi < \lambda$ ”.  $\square$

Readily from the definitions,  $H_0 = \text{Id}$  (the identity function). By theorem 16, 18, and 19,  $H_\sigma$  as defined here agrees with  $H^\sigma$  as defined in [7]. It may be seen that this continues to hold at  $\kappa^+$ ; further discussion is omitted here.

### 8. Enforceability

Let  $\Phi$  be the sentence  $\forall A, C, X, Y, Z$  (“ $A$  is an IV scheme term”  $\wedge$  “ $C$  is club”  $\wedge X = \text{Inac} \wedge Y = H_A(X) \wedge Z = Y \cap C \Rightarrow Z \neq \emptyset$ ). Most of the subformulas of the matrix are readily seen to be  $\Delta_0^1$ ; discussions will be given for the remaining ones.

There is a  $\Delta_0^1$  predicate stating that the class  $X$  represents a well-order on a subset of  $\kappa$ . This states that  $X$  is a class of ordered pairs, which as a binary relation is transitive and reflexive, and has no descending chains of length  $\omega$ .

There is a  $\Delta_0^1$  predicate stating that the class  $X$  represents a scheme for  $\kappa$ . Namely, it represents a pair  $\langle \sigma, \phi \rangle$  where  $\sigma$  is represented as above, and  $\phi$  is a function whose domain is the limit points  $\alpha < \sigma$ , where  $\phi(\alpha)$  is a function with domain either an ordinal, or all ordinals, etc.

An IV term may be given as a class coding the sequence of classes  $\langle t, \sigma_1, \dots, \sigma_k \rangle$  where  $t$  is a hereditarily finite set coding the tree of the term, and  $\sigma_1, \dots, \sigma_k$  are (codes for) the schemes at the leaves in order. It follows that the formula “ $A$  is an IV scheme term” is  $\Delta_0^1$ .

A function  $F : \text{Ord} \mapsto \text{Ord}$  is a class (namely its ordered pairs). Let  $f_x$  denote the value of  $f_\sigma$  at stage  $x$  of the iteration. The class of triples  $\langle x, \lambda, f_x(\lambda) \rangle$  may serve as a witness that  $F = f_A$  for a scheme  $A$ . The predicate “ $W$  is the witness to  $F = f_A$ ” is  $\Delta_0^1$ .

The predicate  $\alpha = C_{\lambda k}(\eta_1, \sigma_1, \dots, \eta_k, \sigma_k)$ , in the values  $\lambda, k, \langle \eta_1, \sigma_1, \dots, \eta_k, \sigma_k \rangle$ , is first-order definable, indeed no doubt  $\Delta_1^0$ . The same is true of the predicate  $\alpha = \phi_{\lambda k}(\zeta, \xi_1, \alpha_1, \dots, \xi_k, \alpha_k)$ , Further discussion is omitted here.

The witness to  $F = f_A$  where  $A$  is any term is the sequence of classes  $\langle F_1, \dots, F_r, W_1, \dots, W_s \rangle$  where  $r$  is the number of nodes in the tree of  $A$ ,  $s$  is the number of leaves, for each leaf index  $i$   $W_i$  witnesses that  $F_{l_i} = f_{\sigma_i}$ , and for each interior node index  $i$   $W_i$  witnesses that  $F_{n_i}$  is obtained by applying the appropriate function per  $\lambda$  to the  $F_j$  of the sons. The predicate “ $W$  is the witness to  $F = f_A$ ” is  $\Delta_0^1$ .

From the foregoing the predicate  $F = f_A$  is  $\Sigma_1^1$ , whence  $\Delta_1^1$ .

A term  $\beta$  over  $\lambda \in \text{In}_\kappa$  is a set. The predicates “ $\beta$  is the ordinal of  $\beta$ ” and “ $\gamma = f_\beta(\mu)$ ” are first order definable. The function  $\lambda \mapsto S_\lambda$  where  $S_\lambda = \{\langle \beta, x, y \rangle : y = H_\beta(x)\}$  may be defined by recursion on  $\lambda$ ; Indeed,  $\langle \beta, x, y \rangle \in S_\lambda$  iff  $\forall \mu(\mu \in y$  iff  $\lambda \in x \wedge \forall \beta(\beta \text{ is term over } \mu \wedge \gamma < f_\beta(\mu) \Rightarrow \text{“}H_\gamma(X \cap \mu)\text{ is stationary”})$ ).

Now,  $Y = H_A(X)$  may be expressed as  $\forall \lambda(\lambda \in Y$  iff  $\lambda \in X \wedge \forall \beta(\beta \text{ is term over } \lambda \wedge \exists \gamma(\langle \lambda, \gamma \rangle \in F \wedge \beta < \gamma \Rightarrow \text{“}H_\beta(X \cap \lambda)\text{ is stationary”})$ ). From the foregoing, “ $Y = H_A(X)$ ” is  $\Delta_1^1$ .

Say that an inaccessible cardinal  $\kappa$  is  $\Lambda_0$ -Mahlo iff  $\models_{V_\kappa} \Phi$ .

**Theorem 21.** *There is a  $\Pi_1^1$  sentence  $\Phi$  such that  $\kappa$  is  $\Lambda_0$ -Mahlo iff  $\models_{V_\kappa} \Phi$ .*

*Proof.* Immediate by remarks above. □

**Lemma 22.** *Suppose  $\kappa$  is weakly compact. For each scheme  $\sigma$  there is a  $\Pi_1^1$  formula  $\Phi_\sigma$  such that  $\models_{V_\kappa} \Phi_\sigma$ , and for  $\lambda \in \text{In}_\kappa$ , if  $\models_{V_\lambda} \Phi_\sigma$  then  $\sigma \cap V_\lambda$  equals  $\sigma \downarrow \lambda$ .*

*Proof.*  $\Phi_\sigma$  is constructed by recursion on  $\sigma$ . Various details are omitted, and may be found in [3]. Cases of the recursion are as follows, where  $\Phi$  has second order parameters  $\alpha, P$ . In case 0,  $\Phi$  is a sentence enforcing that  $\kappa$  is inaccessible. In case 1,  $\Phi_\sigma$  is “ $\sigma$  is a scheme”  $\wedge$  “ $\sigma$  is a successor ordinal”  $\wedge \forall \tau(\tau = \sigma^- \Rightarrow \Phi_{\sigma^-}(\tau, P))$ . In case 2,  $\Phi_\sigma$  is “ $\sigma$  is a scheme”  $\wedge \sigma \in \text{Lim} \wedge \exists \eta(\eta$  is the domain of the ascending sequence of the last node”  $\wedge \forall \xi < \eta \forall \tau(\tau = \sigma_\xi \Rightarrow \Phi_{\sigma_\xi}(\tau, P))$ ). The formulas  $\Phi_{\sigma_\xi}$  may be combined into a single formula in a well-known manner. Case 3 is similar. □

**Corollary 23.** *Suppose  $\kappa$  is weakly compact. For each IV term  $\alpha$  there is a  $\Pi_1^1$  formula  $\Phi_\alpha$  such that  $\models_{V_\kappa} \Phi_\alpha$ , and for  $\lambda \in \text{In}_\kappa$ , if  $\models_{V_\lambda} \Phi_\alpha$  then  $\alpha \cap V_\lambda$  equals  $\alpha \downarrow \lambda$ .*

*Proof.* This follows readily from the lemma and the definition of  $\alpha \downarrow \lambda$ . □

**Theorem 24.** *Suppose  $\kappa$  is weakly compact. Then  $\models_{V_\kappa} \Phi$ .*

*Proof.* It suffices to show by induction on  $\alpha$  that for all  $\alpha$ ,  $\models_{V_\kappa}$  “ $H_\alpha(\text{Inac})$  is stationary”. Inductively,  $\models_{V_\kappa} \forall \beta < \alpha$  “ $H_\beta(\text{Inac})$  is stationary”. By lemma 11 and corollary 23, for a stationary set of  $\lambda$ ,  $\models_{V_\lambda} \forall \beta < f_\alpha(\lambda)$  “ $H_\beta(\text{Inac})$  is stationary”.  $\models_{V_\kappa}$  “ $H_\alpha(\text{Inac})$  is stationary” follows. □

## 9. A New Axiom

It seems helpful to define axiom M to be the second order axiom, to be added to NBG, which states that Ord is "...-Mahlo" where "..." is a specification of the length of stationary set chains which have been constructed. As of this paper, axiom M states that "Ord is  $\Lambda_0$ -Mahlo". The sentence  $\Phi$  of the previous section gives a more detailed statement.

Only some brief remarks on justifying this axiom will be given here. The reader is assumed to be familiar with section 7 of [6]. Inductively, it may be assumed that  $\forall \beta < \alpha$  " $H_\beta(\text{Inac})$ ". Indeed, letting  $\alpha$  be a term in some sufficiently large outer universe, successive  $\lambda$  such that  $\models_{V_\lambda} \forall \beta < f_\alpha(\lambda)$  " $H_\beta(\text{Inac})$  is stationary" may be collected. The general fact that something that may be repeated may be repeated stationarily often results in a universe in which " $H_\alpha(\text{Inac})$  is stationary" holds.

As usual, more details of this justification would be desirable, but we omit this subject here.

## 10. Normal Ultrafilters

An application of function chains was given in [5], with further discussion given in [7]. An essential fact is the following.

**Theorem 25.** *For any measurable cardinal  $\kappa$ , any normal ultrafilter  $U$  on  $\kappa$ , and any IV scheme term  $\alpha$  over  $\kappa$ ,  $f_\alpha$  represents  $\alpha$  in the ultrapower  $Ult_U(V)$ .*

*Proof.* The proof is by induction on  $\alpha$ , using Los' theorem. First the claim is proved for schemes. For case 0 the claim is immediate. For case 1, inductively  $f_\tau$  represents  $\tau$ , whence  $f_{\tau+1}$  represents  $\tau+1$ . For case 2, inductively  $f_{\sigma_\xi}$  represents  $\sigma_\xi$ , whence  $\sup_\xi f_{\sigma_\xi}$  represents  $\sup_\xi \sigma_\xi$ . Case 3 is similar. For the induction on scheme terms, cases 0 and 1 have already been proved. Cases 2 and 3 follow readily by the definability of the functions  $C_k$  and  $\phi_k$ .  $\square$

It is independent whether there are canonical functions of rank  $\kappa^+$  (see [1]). The foregoing shows that the functions  $f_\alpha$  have a weaker property of interest. It would be of interest to obtain a characterization of the property which made no use of ultrafilters.

Suppose  $\kappa$  is a measurable cardinal. Let  $o(\kappa)$  denote its Mitchell order. For  $\alpha$  an IV scheme term let  $S_\alpha$  denote  $\{\lambda \in \text{In}_\kappa : o(\lambda) \geq f_\alpha(\lambda)\}$ . Recall the definition of  $<_R$  from [7].

**Theorem 26.** *Suppose  $\alpha < o(\kappa)$ . Then  $S_\alpha$  is stationary, and if  $\beta < \alpha$  then  $S_\beta <_R S_\alpha$ .*

*Proof.* Let  $U_1$  be a normal ultrafilter on  $\kappa$  with  $O(U_1) = \alpha$ . By lemma 19.34 of [Jech],  $o$  represents  $o(U_1)$  in  $\text{Ult}_{U_1}(V)$ . Since by theorem 25  $f_\alpha$  also represents  $\alpha = o(U_1)$ ,  $\{\lambda \in \text{In}_\kappa : o(\lambda) = f_\alpha(\lambda)\} \in U_1$ . It follows that  $S_\alpha$  is stationary.

Suppose  $\lambda \in S_{\beta+1}$ . Then  $o(\lambda) \geq f_{\beta+1}(\lambda)$ , so except for a thin set of  $\lambda$  there is a normal ultrafilter  $U'$  on  $\lambda$  with  $o(U') = f_\beta(\lambda)$ . By lemma 11 and theorem 25, except for a thin set of  $\lambda$ ,  $f_\beta \upharpoonright \lambda$  represents  $f_\beta(\lambda)$  in  $\text{Ult}_{U'}(V)$ , whence it represents  $o(U')$ . By an argument just given,  $\{\mu < \lambda : o(\mu) = f_\alpha(\mu)\} \in U'$ , whence  $S_\beta \cap \lambda \in U'$ . This shows that  $S_{\beta+1} \subseteq_t H(S_\beta)$ , completing the proof of the theorem.  $\square$

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