

## ON STRICT-DOUBLE-BOUND NUMBERS OF COMPLETE GRAPHS WITHOUT EDGES OF CYCLES

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**Abstract:** For a poset  $P = (X, \leq_P)$ , the *strict-double-bound graph* of  $P$  is the graph  $\text{sDB}(P)$  on  $V(\text{sDB}(P)) = X$  for which vertices  $u$  and  $v$  of  $\text{sDB}(P)$  are adjacent if and only if  $u \neq v$  and there exist elements  $x, y \in X$  distinct from  $u$  and  $v$  such that  $x \leq_P u \leq_P y$  and  $x \leq_P v \leq_P y$ . The *strict-double-bound number*  $\zeta(G)$  of a graph  $G$  is defined as  $\min\{n; \text{sDB}(P) \cong G \cup \overline{K}_n \text{ for some poset } P\}$ . We obtain upper bounds of strict-double-bound numbers of  $K_n - E(C_m)$  ( $5 \leq m \leq n$ ).

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### 1. Introduction

In this paper we consider graphs without loops and multiple edges. For a graph  $G$  and  $S \subseteq V(G)$ ,  $\langle S \rangle_V$  is the induced subgraph of  $S$ . For a graph  $G$  and  $v \in V(G)$ ,  $N_G(v) = \{u; uv \in E(G)\}$ . For a graph  $G$ ,  $\overline{G}$  is the complement of  $G$ . For a graph  $G$  and a subgraph  $H$  of  $G$ , the graph  $G - E(H)$  is the graph with the vertex set  $V(G - E(H)) = V(G)$  and the edge set  $E(G - E(H)) = E(G) - E(H)$ .

The *union*  $G \cup I$  of two graphs  $G$  and  $I$  is the graph with the vertex set  $V(G \cup I) = V(G) \cup V(I)$  and the edge set  $E(G \cup I) = E(G) \cup E(I)$ . The *sum*  $G + I$  of two graphs  $G$  and  $I$  is the graph with the vertex set  $V(G + I) = V(G) \cup V(I)$  and the edge set  $E(G + I) = E(G) \cup E(I) \cup \{uv; u \in V(G), v \in V(I)\}$ .

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A *clique* in a graph  $G$  is the vertex set of a maximal complete subgraph of  $G$ . A family  $\mathcal{Q} = \{Q_1, Q_2, \dots, Q_l\}$  is an *edge clique cover* of  $G$  if each  $Q_i$  is a clique of  $G$ , and for each  $uv \in E(G)$ , there exists  $Q_i \in \mathcal{Q}$  such that  $u, v \in Q_i$ .

For a poset  $P$ , let  $\text{Max}(P)$  be the set of all maximal elements of  $P$  and  $\text{Min}(P)$  be the set of all minimal elements of  $P$ . For a poset  $P$  and elements  $u$  and  $v$ ,  $u \parallel v$  denotes that  $u$  is incomparable with  $v$ .

We consider strict-double-bound graphs and strict-double-bound numbers. For a poset  $P = (X, \leq_P)$ , the *strict-double-bound graph* (*sDB-graph*) of  $P$  is the graph  $\text{sDB}(P)$  on  $V(\text{sDB}(P)) = X$  for which vertices  $u$  and  $v$  of  $\text{sDB}(P)$  are adjacent if and only if  $u \neq v$  and there exist elements  $x, y \in X$  distinct from  $u$  and  $v$  such that  $x \leq_P u \leq_P y$  and  $x \leq_P v \leq_P y$ . We say that a graph  $G$  is a *strict-double-bound graph* if there exists a poset whose strict-double-bound graph is isomorphic to  $G$ . McMorris and Zaslavsky [5] introduced a concept of strict-double-bound graphs. Note that maximal elements and minimal elements of a poset  $P$  are isolated vertices of  $\text{sDB}(P)$ . So a connected graph with  $p \geq 2$  vertices is not a strict-double-bound graph. Scott [7] showed as follows.

**Theorem 1.1** (Scott [7]). *Any graph that is the disjoint union of a non-trivial component and enough number of isolated vertices is a strict-double-bound graph.*

We introduce the strict-double-bound number of a graph. The *strict-double-bound number*  $\zeta(G)$  of a graph  $G$  is defined as  $\min\{n ; \text{sDB}(P) \cong G \cup \overline{K}_n \text{ for some poset } P\}$ .

Scott [7] obtained the following result, using a concept of transitive double competition numbers.

**Theorem 1.2** (Scott [7]). *For a non-trivial connected graph  $G$  and a minimal edge clique cover  $\mathcal{Q}$  of  $G$ ,  $\lceil 2\sqrt{|\mathcal{Q}|} \rceil \leq \zeta(G) \leq |\mathcal{Q}| + 1$ .*

We already knew the following results.

**Proposition 1.3** (Kanada, Kushima, Ogawa, Tagusari and Tsuchiya [2]). *Let  $G$  be a connected graph with  $p \geq 2$  vertices and  $P$  a poset with  $\text{sDB}(P) \cong G \cup \overline{K}_{\zeta(G)}$ . Then  $|\text{Max}(P) \cup \text{Min}(P)| = \zeta(G)$ .*

**Theorem 1.4** (Konishi, Ogawa, Tagusari and Tsuchiya [4]). *For a graph  $G$  with  $p \geq 2$  vertices and no isolated vertices,  $\zeta(K_n + G) = \zeta(G)$  for  $n \geq 1$ .*

Using these results, we deal with a strict-double-bound number of  $G - E(H)$  for some graph  $H$ . For example,  $H$  is a graph with  $|E(G)| \leq 4$  ([1], [6]). And  $H$  is  $K_n$ ,  $K_{1,n}$ ,  $K_n$ -pan ([2]). In this paper we consider  $K_n - E(C_m)$ .

## 2. Cycles

Kim, Park and Sano [3] gave the following result. For a graph  $G$ ,  $\theta_e(G)$  is the edge clique number of  $G$ , that is,  $\theta_e(G) = \min \{|\mathcal{Q}|; \mathcal{Q} \text{ is an edge clique cover of } G\}$ .

**Theorem 2.1** (Kim, Park and Sano [3]). (1)  $\theta_e(\overline{C_6}) = 5$ ,

(2)  $\theta_e(\overline{C_7}) = 7$ ,

(3)  $\theta_e(\overline{C_8}) = 6$ ,

(4)  $\theta_e(\overline{C_m}) \leq \frac{m+5}{2}$  ( $m \geq 9, m : \text{odd}$ ),

(5)  $\theta_e(\overline{C_m}) \leq \frac{m+2}{2}$  ( $m \geq 10, m : \text{even}$ ).

Konishi, Ogawa, Tagusari and Tsuchiya [4] obtained the following result.

**Proposition 2.2** (Konishi, Ogawa, Tagusari and Tsuchiya [4]). For a cycle  $C_n$  with  $n \geq 4$ ,  $\zeta(C_n) = \lfloor 2\sqrt{n} \rfloor$ .

We obtain the following results. For a poset  $P$  and  $v \in V(P)$ ,  $\text{Max}_P(v) = \{u \in \text{Max}(P); v \leq_P u\}$  and  $\text{Min}_P(v) = \{u \in \text{Min}(P); u \leq_P v\}$ .

**Lemma 2.3.** For a graph  $K_m - E(C_m)$  and  $m \geq 5$ , let  $P_{K_m - E(C_m)}$  be a poset such that  $\text{sDB}(P_{K_m - E(C_m)}) \cong (K_m - E(C_m)) \cup \overline{K}_{\zeta(K_m - E(C_m))}$ .

(1) For  $\forall v \in V(K_m - E(C_m))$ ,  $|\text{Max}_{P_{K_m - E(C_m)}}(v)| \geq 2$  or  $|\text{Min}_{P_{K_m - E(C_m)}}(v)| \geq 2$ .

(2) There exist no element  $v$  such that  $\text{Max}_{P_{K_m - E(C_m)}}(v) = \text{Max}(P_{K_m - E(C_m)})$  and  $\text{Min}_{P_{K_m - E(C_m)}}(v) = \text{Min}(P_{K_m - E(C_m)})$ .

*Proof.* (1) We assume that there exists  $v \in V(K_m - E(C_m))$  such that  $\text{Max}_{P_{K_m - E(C_m)}}(v) = \{\alpha\}$  and  $\text{Min}_{P_{K_m - E(C_m)}}(v) = \{\beta\}$ . Then  $\beta \leq_{P_{K_m - E(C_m)}} \alpha$  for  $\forall u \in N_{K_m - E(C_m)}(v)$ . Thus  $N_{K_m - E(C_m)}(v)$  is a clique of  $K_m - E(C_m)$ . Since  $m \geq 5$ , we have a contradiction.

(2) We assume that there exists  $v \in V(K_m - E(C_m))$  such that

$$\text{Max}_{P_{K_m - E(C_m)}}(v) = \text{Max}(P_{K_m - E(C_m)})$$

and

$$\text{Min}_{P_{K_m - E(C_m)}}(v) = \text{Min}(P_{K_m - E(C_m)}).$$

Thus

$$N_{K_m - E(C_m)}(v) = V(K_m - E(C_m)) - \{v\},$$

which is a contradiction.

Based on these results, we obtain the following results.

**Proposition 2.4.** (i)  $\zeta(K_n - E(C_5)) = 5$  ( $n \geq 5$ ),

(ii)  $\zeta(K_n - E(C_6)) = 6$  ( $n \geq 6$ ),

(iii)  $\zeta(K_n - E(C_7)) = 7$  ( $n \geq 7$ ),

(iv)  $\zeta(K_n - E(C_8)) = 7$  ( $n \geq 8$ ),

(v)  $\zeta(K_n - E(C_m)) \leq \frac{m+5}{2} + 1$  ( $n \geq m \geq 9$ ,  $m$  : odd),

(vi)  $\zeta(K_n - E(C_m)) \leq \frac{m+2}{2} + 1$  ( $n \geq m \geq 10$ ,  $m$  : even).

*Proof.* The graph  $K_m - E(C_m)$  is  $\overline{C}_m$  and  $K_m - E(C_m)$  is a connected graph if  $m \geq 5$ . And  $K_n - E(C_m) = K_{n-m} + (K_m - E(C_m)) = K_{n-m} + \overline{C}_m$ . By Theorem 1.4,  $\zeta(K_n - E(C_m)) = \zeta(K_m - E(C_m)) = \zeta(\overline{C}_m)$ .

(i) Since  $\overline{C}_5 = C_5$ ,  $\zeta(C_5) = 5$  by Proposition 2.2. Thus  $\zeta(K_n - E(C_5)) = 5$  ( $n \geq 5$ ).

By Theorem 1.2,  $\left\lceil 2\sqrt{\theta_e(\overline{C}_m)} \right\rceil \leq \zeta(\overline{C}_m) \leq \theta_e(\overline{C}_m) + 1$ . Using Theorem 2.1, we have the followings:

(ii)  $5 \leq \zeta(K_n - E(C_6)) \leq 6$  ( $n \geq 6$ ),

(iii)  $6 \leq \zeta(K_n - E(C_7)) \leq 8$  ( $n \geq 7$ ),

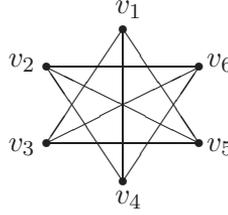
(iv)  $5 \leq \zeta(K_n - E(C_8)) \leq 7$  ( $n \geq 8$ ),

(v)  $\zeta(K_n - E(C_m)) \leq \frac{m+5}{2} + 1$  ( $n \geq m \geq 9$ ,  $m$  : odd),

(vi)  $\zeta(K_n - E(C_m)) \leq \frac{m+2}{2} + 1$  ( $n \geq m \geq 10$ ,  $m$  : even).

(ii)  $\zeta(K_n - E(C_6)) = 6$  ( $n \geq 6$ ) :

We consider a graph  $K_6 - E(C_6)$  (see Figure 1). We assume that there exists a poset  $P$  such that  $\text{sDB}(P) \cong (K_6 - E(C_6)) \cup \overline{K}_5$ . In the case of  $|\text{Max}(P)| = 4$

Figure 1:  $K_6 - E(C_6)$ .

and  $|\text{Min}(P)| = 1$  (or  $|\text{Max}(P)| = 1$  and  $|\text{Min}(P)| = 4$ )  $\text{sDB}(P)$  has at most four cliques. So we consider the case  $|\text{Max}(P)| = 3$  and  $|\text{Min}(P)| = 2$  by Theorem 2.1. And the case  $|\text{Max}(P)| = 2$  and  $|\text{Min}(P)| = 3$  is the dual. Let  $\text{Max}(P) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}(P) = \{\beta_1, \beta_2\}$ .

**Case (ii)-1 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

Then for adjacent vertices of  $v_1, v_3, v_4, v_5 \leq_P \alpha_1, \beta_1 \leq_P v_3, v_5$ , and  $\beta_2 \leq_P v_4$ , because  $v_4$  is not adjacent to  $v_3$  and  $v_5$ . Since  $N_G(v_2) = \{v_4, v_5, v_6\}$  and  $N_G(v_6) = \{v_2, v_3, v_4\}$ ,  $\beta_1 \leq_P v_2, v_6$  and  $\beta_2 \leq_P v_2, v_6$ . And also  $v_2, v_5 \leq_P \alpha_2$  and  $v_6, v_3 \leq_P \alpha_3$ . Since  $v_2$  is adjacent to  $v_6$ ,  $v_2 \leq_P \alpha_3$  or  $v_6 \leq_P \alpha_2$ . Then  $v_2$  is adjacent to  $v_3$  or  $v_6$  is adjacent to  $v_5$ , which is a contradiction.

**Case (ii)-2 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

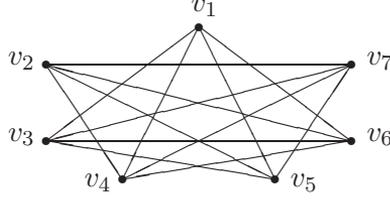
Since  $v_2$  is not adjacent to  $v_1$ ,  $\text{Max}_P(v_2) = \{\alpha_3\}$ . Then we have Case (ii)-1 or  $|\text{Max}_P(v_2)| = |\text{Min}_P(v_2)| = 1$ .

**Case (ii)-3 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

Then for adjacent vertices of  $v_1, \beta_1 \leq_P v_3, v_4, v_5$ , and  $v_3, v_5 \leq_P \alpha_1$ , and  $v_4 \leq_P \alpha_2$ , because  $v_4$  is not adjacent to  $v_3$  and  $v_5$ . Since  $v_3$  is not adjacent to  $v_4$ ,  $v_3 \parallel \alpha_3$  or  $v_4 \parallel \alpha_3$ . Thus  $|\text{Max}_P(v_3)| = 1$  or  $|\text{Max}_P(v_4)| = 1$ . So we have Case (ii)-1,  $|\text{Max}_P(v_3)| = |\text{Min}_P(v_3)| = 1$  or  $|\text{Max}_P(v_4)| = |\text{Min}_P(v_4)| = 1$ .

**Case (ii)-4 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

Then for adjacent vertices of  $v_1, \beta_1 \leq_P v_3, v_4, v_5$ . Since  $v_3$  is not adjacent to  $v_4$ ,  $|\text{Max}_P(v_3)| = 1$  or  $|\text{Max}_P(v_4)| = 1$ . Then we have Case (ii)-1,

Figure 2:  $K_7 - E(C_7)$ .

$$|\text{Max}_P(v_3)| = |\text{Min}_P(v_3)| = 1 \text{ or } |\text{Max}_P(v_4)| = |\text{Min}_P(v_4)| = 1.$$

Thus  $\zeta(K_6 - E(C_6)) \geq 6$ , and  $\zeta(K_n - E(C_6)) = \zeta(K_6 - E(C_6)) = 6$ .

(iii)  $\zeta(K_n - E(C_7)) = 7$  ( $n \geq 7$ ):

We consider a graph  $K_7 - E(C_7)$  (see Figure 2). We assume that there exists a poset  $P$  such that  $\text{sDB}(P) \cong (K_7 - E(C_7)) \cup \overline{K_6}$ . In the case of  $|\text{Max}(P)| = 5$  and  $|\text{Min}(P)| = 1$  (or  $|\text{Max}(P)| = 1$  and  $|\text{Min}(P)| = 5$ ),  $\text{sDB}(P)$  has at most five cliques. So we consider the case  $|\text{Max}(P)| = 4$  and  $|\text{Min}(P)| = 2$ , and the case  $|\text{Max}(P)| = |\text{Min}(P)| = 3$  by Theorem 2.1. And the case  $|\text{Max}(P)| = 2$  and  $|\text{Min}(P)| = 4$  is the dual. First we consider the case  $|\text{Max}(P)| = 4$  and  $|\text{Min}(P)| = 2$ . Let  $\text{Max}(P) = \{\alpha_1, \alpha_2, \alpha_3, \alpha_4\}$  and  $\text{Min}(P) = \{\beta_1, \beta_2\}$ .

**Case (iii)-1 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

Then for adjacent vertices of  $v_1, v_3, v_4, v_5, v_6 \leq_P \alpha_1$ . Since  $v_4$  is not adjacent to  $v_3$  and  $v_5$ ,  $\beta_1 \leq_P v_3, v_5$ , and  $\beta_2 \leq_P v_4$ . Since  $v_6$  is adjacent to  $v_3$ ,  $\beta_1 \leq_P v_6$ . Thus  $v_6$  is adjacent to  $v_5$ , which is a contradiction.

**Case (iii)-2 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

Then for adjacent vertices of  $v_1, \beta_1 \leq_P v_3, v_4, v_5, v_6$ . Since  $v_4$  is not adjacent to  $v_3$  and  $v_5$ ,  $v_3, v_5 \leq_P \alpha_1$ , and  $v_4 \leq_P \alpha_2$ . Since  $v_6$  is adjacent to  $v_1, v_3, v_4$  and not adjacent to  $v_5$ ,  $v_6 \leq_P \alpha_2$  and  $v_6, v_3 \leq_P \alpha_3$ .

Since  $v_4$  is not adjacent to  $v_5$ ,  $\text{Max}_P(v_4) = \{\alpha_2\}$  or  $\text{Max}_P(v_5) = \{\alpha_1\}$ . Then we have Case (iii)-1,  $|\text{Max}_P(v_4)| = |\text{Min}_P(v_4)| = 1$  or  $|\text{Max}_P(v_5)| = |\text{Min}_P(v_5)| = 1$ .

**Case (iii)-3 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

Then for adjacent vertices of  $v_1$ ,  $\beta_1 \leq_P v_3, v_4, v_5, v_6$ . Since  $\langle\{v_5, v_3, v_6, v_4\}\rangle_V$  is a path, we have two cases: (1)  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_2$  and  $v_6, v_4 \leq_P \alpha_3$ , (2)  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_4$  and  $v_6, v_4 \leq_P \alpha_3$ .

In the case  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_2$  and  $v_6, v_4 \leq_P \alpha_3$ , if  $\text{Max}_P(v_4) = \{\alpha_3\}$  or  $\text{Max}_P(v_5) = \{\alpha_1\}$ , we have Case (iii)-1,  $|\text{Max}(v_4)| = |\text{Min}(v_4)| = 1$  or  $|\text{Max}(v_5)| = |\text{Min}(v_5)| = 1$ . So  $v_4, v_5 \leq_P \alpha_4$ . Then  $v_4$  is adjacent to  $v_5$ , which is a contradiction.

In the case  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_4$  and  $v_6, v_4 \leq_P \alpha_3$ , if  $\text{Max}_P(v_4) = \{\alpha_3\}$  or  $\text{Max}_P(v_5) = \{\alpha_1\}$ , then we have Case (iii)-1,  $|\text{Max}(v_4)| = |\text{Min}(v_4)| = 1$  or  $|\text{Max}(v_5)| = |\text{Min}(v_5)| = 1$ . So  $v_4, v_5 \leq_P \alpha_2$ , because  $v_4$  is not adjacent to  $v_3$  and  $v_5$  is not adjacent to  $v_6$ . Then  $v_4$  is adjacent to  $v_5$ , which is a contradiction.

**Case (iii)-4 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3, \alpha_4\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

Then for adjacent vertices of  $v_1$ ,  $\beta_1 \leq_P v_3, v_4, v_5, v_6$ . Since  $\langle\{v_5, v_3, v_6, v_4\}\rangle_V$  is a path,  $v_5, v_3 \leq_P \alpha_1$ , and  $v_3, v_6 \leq_P \alpha_2$  and  $v_6, v_4 \leq_P \alpha_3$ ,

If  $\text{Max}_P(v_4) = \{\alpha_3\}$  or  $\text{Max}_P(v_5) = \{\alpha_1\}$ , we have Case (iii)-1,  $|\text{Max}(v_4)| = |\text{Min}(v_4)| = 1$  or  $|\text{Max}(v_5)| = |\text{Min}(v_5)| = 1$ . So  $v_4, v_5 \leq_P \alpha_4$ , because  $v_4$  is not adjacent to  $v_3$  and  $v_5$  is not adjacent to  $v_6$ . Then  $v_4$  is adjacent to  $v_5$ , which is a contradiction.

**Case (iii)-5 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

If  $|\text{Min}_P(v_i)| = 1$  ( $i = 2, 3, \dots, 7$ ), we have Case (iii)-2, Case (iii)-3, Case (iii)-4, or  $|\text{Max}_P(v_i)| = |\text{Min}_P(v_i)| = 1$ . We consider the case  $\beta_1 \leq_P v_2, v_3, v_4, v_5, v_6, v_7$  and  $\beta_2 \leq_P v_2, v_3, v_4, v_5, v_6, v_7$ .

Then  $\text{sDB}(P)$  has at most four cliques,  $C_i = \{w \notin \text{Max}(P) \cup \text{Min}(P) ; w \leq_P \alpha_i\}$  ( $i = 1, 2, 3, 4$ ), which is a contradiction.

**Case (iii)-6 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

Since  $v_2$  is not adjacent to  $v_1$ ,  $v_2 \leq_P \alpha_4$ . We have Case (iii)-1 or  $|\text{Max}_P(v_2)| = |\text{Min}_P(v_2)| = 1$ .

Next we consider the case  $|\text{Max}(P)| = 3$  and  $|\text{Min}(P)| = 3$ . Let  $\text{Max}(P) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}(P) = \{\beta_1, \beta_2, \beta_3\}$ .

**Case (iii)-7 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

Then for adjacent vertices of  $v_1, v_3, v_4, v_5, v_6 \leq_P \alpha_1$ . Since  $v_4$  is not adjacent to  $v_3$  and  $v_5$ ,  $\beta_1 \leq_P v_3, v_5$  and  $\beta_2 \leq_P v_4$ . Since  $v_6$  is adjacent to  $v_1$  and  $v_3$ , we have two cases: (1)  $\beta_1 \leq_P v_6$ , and (2)  $\beta_2 \leq_P v_6$  and  $\beta_3 \leq_P v_3, v_6$ .

In the case  $\beta_1 \leq_P v_6$ ,  $v_6$  is adjacent to  $v_5$ . We have a contradiction.

In the case  $\beta_2 \leq_P v_6$  and  $\beta_3 \leq_P v_3, v_6$ , the element  $v_4$  is incomparable with  $\beta_1, \beta_3$ , and  $v_5$  is incomparable with  $\beta_2, \beta_3$ . Since  $v_2$  is adjacent to  $v_4$  and  $v_5$ ,  $\beta_1, \beta_2 \leq_P v_2$ .

Since  $v_7$  is adjacent to  $v_4$  and  $v_5$ ,  $\beta_1, \beta_2 \leq_P v_7$ . And  $v_2 \parallel \alpha_1$  and  $v_7 \parallel \alpha_1$ , because  $v_2$  and  $v_7$  are not adjacent to  $v_1$

If  $v_2 \leq_P \alpha_2$  and  $v_2 \parallel \alpha_3$ , then  $v_4, v_5, v_6, v_7 \leq_P \alpha_2$ , because  $v_2$  is adjacent to  $v_4, v_5, v_6$  and  $v_7$ . Then  $v_7$  is adjacent to  $v_6$ , which is a contradiction. So  $v_2 \leq_P \alpha_2$  and  $v_2 \leq_P \alpha_3$ . Since  $v_7$  is adjacent to  $v_3, v_7, v_3 \leq_P \alpha_2$  or  $v_7, v_3 \leq_P \alpha_3$ . Then  $v_2$  is adjacent to  $v_3$ , which is a contradiction.

**Case (iii)-8 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2, \beta_3\}$ .

Then for adjacent vertices of  $v_1, v_3, v_4, v_5, v_6 \leq_P \alpha_1$ . Since  $\langle \{v_5, v_3, v_6, v_4\} \rangle_V$  is a path,  $\beta_1 \leq_P v_5, v_3, \beta_2 \leq_P v_3, v_6$  and  $\beta_3 \leq_P v_6, v_4$ .

Since  $v_2$  and  $v_7$  are adjacent to  $v_4$  and  $v_5$ ,  $\beta_1, \beta_3 \leq_P v_2$  and  $\beta_1, \beta_3 \leq_P v_7$ . And  $v_2 \parallel \alpha_1$  and  $v_7 \parallel \alpha_1$ , because  $v_2$  and  $v_7$  are not adjacent to  $v_1$ . If  $v_2 \leq_P \alpha_2$  and  $v_2 \parallel \alpha_3$ , then  $v_4, v_5, v_6, v_7 \leq_P \alpha_2$ , because  $v_2$  is adjacent to  $v_4, v_5, v_6$  and  $v_7$ . Then  $v_6$  is adjacent to  $v_7$ , which is a contradiction. So  $v_2 \leq_P \alpha_2$  and  $v_2 \leq \alpha_3$ . Since  $v_7$  is adjacent to  $v_3, v_7, v_3 \leq_P \alpha_2$  or  $v_7, v_3 \leq_P \alpha_3$ . Then  $v_2$  is adjacent to  $v_3$ , which is a contradiction.

**Case (iii)-9 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2, \beta_3\}$ .

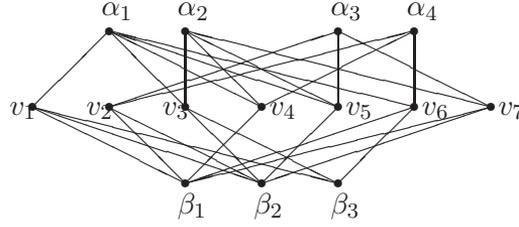
Since  $v_2$  is not adjacent to  $v_1$ ,  $\text{Max}_P(v_2) = \{\alpha_3\}$ . We have Case (iii)-7, Case (iii)-8, or  $|\text{Max}_P(v_2)| = |\text{Min}_P(v_2)| = 1$ .

**Case (iii)-10 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

This case is the dual case of Case (iii)-7.

**Case (iii)-11 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

This case is the dual case of Case (iii)-8.


 Figure 3:  $P_{K_7-E(C_7)}$ .

**Case (iii)-12 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

This case is the dual case of Case (iii)-9.

**Case (iii)-13 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

If  $|\text{Max}_P(v_2)| = 1$ , we have Case (iii)-7, Case (iii)-8, or  $|\text{Min}_P(v_2)| = 1$ . If  $|\text{Min}_P(v_2)| = 1$ , we have Case (iii)-10, Case (iii)-11, or  $|\text{Max}_P(v_2)| = 1$ . Thus  $|\text{Max}_P(v_2)| \geq 2$  and  $|\text{Min}_P(v_2)| \geq 2$ . Then  $v_2$  is adjacent to  $v_1$ , which is a contradiction.

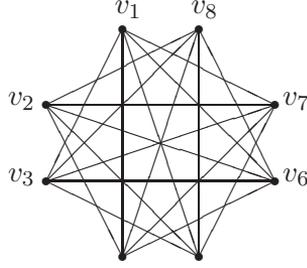
Thus  $\zeta(K_7-E(C_7)) \geq 7$ . We have a poset with  $\text{sDB}(P) \cong (K_7-E(C_7)) \cup \overline{K_7}$  (see Figure 3). Therefore  $\zeta(K_n - E(C_7)) = \zeta(K_7 - E(C_7)) = 7$ .

(iv)  $\zeta(K_n - E(C_8)) = 7$  ( $n \geq 8$ ) :

We consider a graph  $K_8 - E(C_8)$  (see Figure 4). We assume that there exists a poset  $P$  such that  $\text{sDB}(P) \cong (K_8 - E(C_8)) \cup \overline{K_6}$ . In the case of  $|\text{Max}(P)| = 5$  and  $|\text{Min}(P)| = 1$  (or  $|\text{Max}(P)| = 1$  and  $|\text{Min}(P)| = 5$ )  $\text{sDB}(P)$  has at most five cliques. So we consider the case  $|\text{Max}(P)| = 4$  and  $|\text{Min}(P)| = 2$ , and the case  $|\text{Max}(P)| = |\text{Min}(P)| = 3$  by Theorem 2.1. And the case  $|\text{Max}(P)| = 2$  and  $|\text{Min}(P)| = 4$  is the dual. First we consider the case  $|\text{Max}(P)| = 4$  and  $|\text{Min}(P)| = 2$ . Let  $\text{Max}(P) = \{\alpha_1, \alpha_2, \alpha_3, \alpha_4\}$  and  $\text{Min}(P) = \{\beta_1, \beta_2\}$ .

**Case (iv)-1 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

Then for adjacent vertices of  $v_1, v_3, v_4, v_5, v_6 \leq_P \alpha_1$ . Since  $v_4$  is not adjacent to  $v_3$  and  $v_5$ ,  $\beta_1 \leq_P v_3, v_5$ , and  $\beta_2 \leq_P v_4$ . Since  $v_6$  is adjacent to  $v_3$ ,  $\beta_1 \leq_P v_6$ .

Figure 4:  $K_8 - E(C_8)$ .

Then  $v_6$  is adjacent to  $v_5$ , which is a contradiction.

**Case (iv)-2 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

Then for adjacent vertices of  $v_1$ ,  $\beta_1 \leq_P v_3, v_4, v_5, v_6$ . Since  $v_4$  is not adjacent to  $v_3$  and  $v_5$ ,  $v_3, v_5 \leq_P \alpha_1$ , and  $v_4 \leq_P \alpha_2$ . Since  $v_6$  is adjacent to  $v_1, v_3, v_4$  and not adjacent to  $v_5$ ,  $v_6 \leq_P \alpha_2$  and  $v_6, v_3 \leq_P \alpha_3$ .

Since  $v_4$  is not adjacent to  $v_5$ ,  $\text{Max}_P(v_4) = \{\alpha_2\}$  or  $\text{Max}_P(v_5) = \{\alpha_1\}$ . Then we have Case (iv)-1,  $|\text{Max}_P(v_4)| = |\text{Min}_P(v_4)| = 1$  or  $|\text{Max}_P(v_5)| = |\text{Min}_P(v_5)| = 1$ .

**Case (iv)-3 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

Then for adjacent vertices of  $v_1$ ,  $\beta_1 \leq_P v_3, v_4, v_5, v_6$ . Since  $\langle \{v_5, v_3, v_6, v_4\} \rangle_V$  is a path, we have two cases: (1)  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_2$  and  $v_6, v_4 \leq_P \alpha_3$ , (2)  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_4$  and  $v_6, v_4 \leq_P \alpha_3$ .

In the case  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_2$  and  $v_6, v_4 \leq_P \alpha_3$ , if  $\text{Max}_P(v_4) = \{\alpha_3\}$  or  $\text{Max}_P(v_5) = \{\alpha_1\}$ , we have Case (iv)-1,  $|\text{Max}(v_4)| = |\text{Min}(v_4)| = 1$  or  $|\text{Max}(v_5)| = |\text{Min}(v_5)| = 1$ . So  $v_4, v_5 \leq_P \alpha_4$ , because  $v_4$  is not adjacent to  $v_3$  and  $v_5$  is not adjacent to  $v_6$ . Then  $v_4$  is adjacent to  $v_5$ , which is a contradiction.

In the case  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_4$  and  $v_6, v_4 \leq_P \alpha_3$ , if  $\text{Max}_P(v_4) = \{\alpha_3\}$  or  $\text{Max}_P(v_5) = \{\alpha_1\}$ , we have Case (iv)-1,  $|\text{Max}(v_4)| = |\text{Min}(v_4)| = 1$  or  $|\text{Max}(v_5)| = |\text{Min}(v_5)| = 1$ . So  $v_4, v_5 \leq_P \alpha_2$ , because  $v_4$  is not adjacent to  $v_3$  and  $v_5$  is not adjacent to  $v_6$ . Then  $v_4$  is adjacent to  $v_5$ , which is a contradiction.

**Case (iv)-4 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3, \alpha_4\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

Then for adjacent vertices of  $v_1$ ,  $\beta_1 \leq_P v_3, v_4, v_5, v_6$ . Since  $\langle \{v_5, v_3, v_6, v_4\} \rangle_V$  is a path,  $v_5, v_3 \leq_P \alpha_1$ ,  $v_3, v_6 \leq_P \alpha_2$  and  $v_6, v_4 \leq_P \alpha_3$ .

If  $\text{Max}_P(v_4) = \{\alpha_3\}$  or  $\text{Max}_P(v_5) = \{\alpha_1\}$ , we have Case (iv)-1,  $|\text{Max}_P(v_4)| = |\text{Min}_P(v_4)| = 1$  or  $|\text{Max}_P(v_5)| = |\text{Min}_P(v_5)| = 1$ . So  $v_4, v_5 \leq_P \alpha_4$ , because  $v_4$  is not adjacent to  $v_3$  and  $v_5$  is not adjacent to  $v_6$ . Then  $v_4$  is adjacent to  $v_5$ , which is a contradiction.

**Case (iv)-5 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

If  $|\text{Min}_P(v_i)| = 1$  ( $i = 2, 3, \dots, 8$ ), we have Case (iv)-2, Case (iv)-3, Case (iv)-4 or  $|\text{Max}_P(v_i)| = |\text{Min}_P(v_i)| = 1$ . We consider the case  $\beta_1 \leq_P v_2, v_3, v_4, v_5, v_6, v_7, v_8$  and  $\beta_2 \leq_P v_2, v_3, v_4, v_5, v_6, v_7, v_8$ . Then  $\text{sDB}(P)$  has at most four cliques,  $C_i = \{w \notin \text{Max}(P) \cup \text{Min}(P) ; w \leq_P \alpha_i\}$  ( $i = 1, 2, 3, 4$ ), which is a contradiction.

**Case (iv)-6 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

Since  $v_2$  is not adjacent to  $v_1$ ,  $v_2 \leq_P \alpha_4$ . We have Case (iv)-1 or  $|\text{Max}_P(v_2)| = |\text{Min}_P(v_2)| = 1$ .

Next we consider the case  $|\text{Max}(P)| = 3$  and  $|\text{Min}(P)| = 3$ . Let  $\text{Max}(P) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}(P) = \{\beta_1, \beta_2, \beta_3\}$ .

**Case (iv)-7 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

Then for adjacent vertices of  $v_1$ ,  $v_3, v_4, v_5, v_6, v_7 \leq_P \alpha_1$ . Since  $v_4$  is not adjacent to  $v_3$  and  $v_5$ ,  $\beta_1 \leq_P v_3, v_5$  and  $\beta_2 \leq_P v_4$ . Since  $v_6$  is adjacent to  $v_1$  and  $v_3$ , we have two cases: (1)  $\beta_1 \leq_P v_6$ , and (2)  $\beta_2 \leq_P v_6$  and  $\beta_3 \leq_P v_3, v_6$ .

In the case  $\beta_1 \leq_P v_6$ ,  $v_6$  is adjacent to  $v_5$ . We have a contradiction.

In the case  $\beta_2 \leq_P v_6$  and  $\beta_3 \leq_P v_3, v_6$ ,  $\text{Min}_P(v_4) = \{\beta_2\}$ , because  $v_4$  is not adjacent to  $v_3$ . Since  $v_7$  is adjacent to  $v_4$ ,  $\beta_2 \leq_P v_7 \leq_P \alpha_1$ . Then  $v_7$  is adjacent to  $v_6$ , which is a contradiction.

**Case (iv)-8 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2, \beta_3\}$ .

Then for adjacent vertices of  $v_1$ ,  $v_3, v_4, v_5, v_6, v_7 \leq_P \alpha_1$ . Since  $\langle \{v_5, v_3, v_6, v_4\} \rangle_V$  is a path,  $\beta_1 \leq_P v_5, v_3$ ,  $\beta_2 \leq_P v_3, v_6$  and  $\beta_3 \leq_P v_6, v_4$ . Since  $v_4$  is not adjacent to  $v_3$ ,  $\text{Min}_P(v_4) = \{\beta_3\}$ . Since  $v_7$  is adjacent to  $v_4$ ,  $\beta_3 \leq_P v_7 \leq_P \alpha_1$ . Then  $v_7$  is adjacent to  $v_6$ , which is a contradiction.

**Case (iv)-9 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2, \beta_3\}$ .

Since  $v_2$  is not adjacent to  $v_1$ ,  $\text{Max}_P(v_2) = \{\alpha_3\}$ . We have Case (iv)-7, or Case (iv)-8, or  $|\text{Max}_P(v_2)| = |\text{Min}_P(v_2)| = 1$ .

**Case (iv)-10 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

This case is the dual case of Case (iv)-7.

**Case (iv)-11 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1\}$ .

This case is the dual case of Case (iv)-8.

**Case (iv)-12 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2, \alpha_3\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

This case is the dual case of Case (iv)-9.

**Case (iv)-13 :** there exists an element  $v = v_1$  such that  $\text{Max}_P(v_1) = \{\alpha_1, \alpha_2\}$  and  $\text{Min}_P(v_1) = \{\beta_1, \beta_2\}$ .

If  $|\text{Max}_P(v_2)| = 1$ , we have Case (iv)-7, Case (iv)-8, or  $|\text{Min}_P(v_2)| = 1$ . If  $|\text{Min}_P(v_2)| = 1$ , we have Case (iv)-10, Case (iv)-11, or  $|\text{Max}_P(v_2)| = 1$ . Thus  $|\text{Max}_P(v_2)| \geq 2$  and  $|\text{Min}_P(v_2)| \geq 2$ . Then  $v_2$  is adjacent to  $v_1$ , which is a contradiction.

Thus  $\zeta(K_8 - E(C_8)) \geq 7$ , and  $\zeta(K_n - E(C_8)) = \zeta(K_8 - E(C_8)) = 7$ .

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