

SOLUTION METHODS FOR
MIXED VARIATIONAL INEQUALITIES

Barbara Panicucci¹, Massimo Pappalardo², Mauro Passacantando³ §

^{1,2,3}Department of Applied Mathematics
University of Pisa

Via Buonarroti 1/c, Pisa, 56127, ITALY

¹e-mail: panicucc@mail.dm.unipi.it

²e-mail: pappalardo@dma.unipi.it

³e-mail: passacantando@dma.unipi.it

Abstract: In this paper two descent methods with respect to a gap function for solving a class of monotone mixed variational inequalities are proposed. We show that the two algorithms, based on an exact and an Armijo-type line search procedure, respectively, are globally convergent.

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1. Introduction

Let X be a nonempty, closed, and convex subset of \mathbb{R}^n , $F : \mathbb{R}^n \rightarrow \mathbb{R}^n$ a map, and $f : \mathbb{R}^n \rightarrow \mathbb{R}$ a real valued function. The Mixed Variational Inequality (MVI, for short) problem is to find a point $x^* \in X$ such that

$$\langle F(x^*), y - x^* \rangle + f(y) - f(x^*) \geq 0 \quad \forall y \in X, \quad (1)$$

where $\langle \cdot, \cdot \rangle$ denotes the usual inner product in \mathbb{R}^n .

This problem was originally considered by Lescarret [6] and Browder [2] for its applications in mathematical physics. Afterwards, it has been shown that the MVI problem has a large variety of applications in various fields such as mechanics, economics and operation research, see [1], [5], [7], [10] and references

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§Correspondence author

therein. When $f \equiv 0$, the MVI problem reduces to the classical variational inequality problem.

One of the main approaches for solving the classical variational inequality or the MVI consists in minimizing a gap function associated to the problem. Plenty of corresponding descent type methods have been developed for VIs (see e.g. [3] and references therein). Descent type methods have been also presented for the MVI problem (see [4] and [9]). In [4] the author proposed a descent method which utilizes an inexact Armijo type line search procedure. The convergence was proved by assuming the operator F to be strongly monotone.

In this paper, we devise two global convergent descent algorithms (with an exact and an inexact line search procedure, respectively) with respect to a gap function, for solving MVIs with monotone (not necessarily strongly monotone) operator.

In the rest of the paper we consider the following assumptions.

(A1) The set $X \subseteq \mathbb{R}^n$ is nonempty, closed, and convex; the map $F : Y \rightarrow \mathbb{R}^n$ is continuously differentiable, where Y is an open convex set such that $X \subset Y$; the function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ is convex.

(A2) The map F is monotone on Y , i.e. $\langle F(x) - F(y), x - y \rangle \geq 0$ for all $x, y \in Y$.

(A3) The set X is bounded.

2. Gap Functions

In this paper we consider the following function (see [4]):

$$\varphi_\alpha(x) = \max_{y \in X} \left[\langle F(x), x - y \rangle - \frac{\alpha}{2} \|x - y\|_G^2 + f(x) - f(y) \right], \quad (2)$$

where α is a positive parameter, G is a symmetric positive definite matrix, and $\|\cdot\|_G$ denotes the norm in \mathbb{R}^n defined by $\|x\|_G = \sqrt{\langle x, Gx \rangle}$. It is easy to check that for each $x \in X$ the optimization problem (2) has a unique solution which will be denoted by $y_\alpha(x)$.

Under assumption (A1), the function defined in (2) is a gap function for the problem (1), i.e. it is nonnegative on X and the set of zeros coincide with the set of solutions of (1), see [4]. Therefore, if the problem (1) has a solution, it is equivalent to the following constrained optimization problem:

$$\min_{x \in X} \varphi_\alpha(x). \quad (3)$$

We remark that, by definition, the gap function φ_α is nondifferentiable since the function f is so. However, under the assumption (A1), the function φ_α is locally Lipschitz continuous on X and it is directionally derivable with respect to any direction (see [10, Proposition 4.19]). Moreover, the Clarke subdifferential and the directional derivative of φ_α at any $x \in X$ can be determined explicitly.

Note also that problem (3) may have local minima which differ from the global one. Therefore in order to solve (1) with a descent method with respect to φ_α , we need some monotonicity properties of the operator F . In [4] it has been proved that, if F is strongly monotone on Y (i.e. $\langle F(x) - F(y), x - y \rangle \geq \tau \|x - y\|^2$, for all $x, y \in Y$, for some $\tau > 0$) and x is not a solution of (1), then $y_\alpha(x) - x$ is a descent direction at any point x for the gap function φ_α . When F is only monotone, the vector $y_\alpha(x) - x$ is not necessarily a descent direction for φ_α , but it satisfies a condition which will be exploited in the following and that allows constructing descent methods for the gap function φ_α .

Theorem 1. *Let assumptions (A1) – (A2) be fulfilled and let $\alpha > 0$. Then, for each $x \in X$, the vector $y_\alpha(x) - x$ satisfies the following condition:*

$$\varphi'_\alpha(x; y_\alpha(x) - x) \leq -\varphi_\alpha(x) + \frac{\alpha}{2} \|x - y_\alpha(x)\|_G^2 \leq 0, \tag{4}$$

where $\varphi'_\alpha(x; y_\alpha(x) - x)$ denotes the directional derivative of φ_α at x with respect to the direction $y_\alpha(x) - x$.

Proof. Let $x \in X$. From [10, Proposition 4.19] we have

$$\begin{aligned} \varphi'_\alpha(x; y_\alpha(x) - x) &= \langle F(x) - (\nabla F(x)^\top - \alpha G)(y_\alpha(x) - x), y_\alpha(x) - x \rangle \\ &\quad + f'(x; y_\alpha(x) - x) \\ &= \langle F(x), y_\alpha(x) - x \rangle + \alpha \|y_\alpha(x) - x\|_G^2 \\ &\quad - \langle y_\alpha(x) - x, \nabla F(x)(y_\alpha(x) - x) \rangle + f'(x; y_\alpha(x) - x). \end{aligned} \tag{5}$$

By assumption (A2), it follows that the matrix $\nabla F(x)$ is positive semidefinite (see [3, Proposition 2.3.2]) and hence

$$\langle y_\alpha(x) - x, \nabla F(x)(y_\alpha(x) - x) \rangle \geq 0. \tag{6}$$

Moreover, since the function f is convex we have

$$f'(x; y_\alpha(x) - x) \leq f(y_\alpha(x)) - f(x). \tag{7}$$

Therefore, taking into account of (5), (6), (7), and that $\varphi_\alpha(x) \geq \frac{\alpha}{2} \|x - y_\alpha(x)\|_G^2$ (see [4]), we have

$$\begin{aligned} \varphi'_\alpha(x; y_\alpha(x) - x) &\leq \langle F(x), y_\alpha(x) - x \rangle + \alpha \|y_\alpha(x) - x\|_G^2 + f(y_\alpha(x)) - f(x) \\ &= -\varphi_\alpha(x) + \frac{\alpha}{2} \|x - y_\alpha(x)\|_G^2 \leq 0. \quad \square \end{aligned}$$

Algorithm 1

0. (Initial step)
 Let G be a symmetric positive definite matrix and $\eta \in (0, 1)$.
 Let $\{\alpha_k\}$ be a sequence strictly decreasing to 0.
 Choose any $x^0 \in X$ and set $k = 0$.
1. (Stopping criterion)
If $\varphi_{\alpha_k}(x^k) = 0$
 then STOP,
else set $k = k + 1$.
2. (Minimization of φ_{α_k})
 2a. (Initialization)
 Set $i = 0$ and $z^0 = x^{k-1}$.
 2b. **If** $-\varphi_{\alpha_k}(z^i) + \frac{\alpha_k}{2} \|z^i - y_{\alpha_k}(z^i)\|_G^2 < -\eta \varphi_{\alpha_k}(z^i)$
 then (line search)
 set $d^i = y_{\alpha_k}(z^i) - z^i$
 compute $t_i \in \arg \min_{t \in [0,1]} \varphi_{\alpha_k}(z^i + t d^i)$
 else (update of x^k)
 set $x^k = z^i$ and return to step 1.
 2c. (Update of z^i)
 Set $z^{i+1} = z^i + t_i d^i$, $i = i + 1$, and return to step 2b.

This result is useful to derive a modified descent method with exact line search procedure (Section 3) and one with inexact Armijo-type line search procedure (Section 4) for solving the MVI problem (1). The basic idea is to use (4) to obtain, if possible, a descent direction. Indeed, if $x \in X$ satisfies the condition

$$-\varphi_{\alpha}(x) + \frac{\alpha}{2} \|x - y_{\alpha}(x)\|_G^2 < -\eta \varphi_{\alpha}(x), \quad (8)$$

where $\eta \in (0, 1)$, then from (4) and (8) we get

$$\varphi'_{\alpha}(x; y_{\alpha}(x)) < -\eta \varphi_{\alpha}(x).$$

Hence the vector $d = y_{\alpha}(x) - x$ is a descent direction for φ_{α} at x and we can perform a line search procedure with respect to the direction d . Otherwise, if x does not solve the problem (1) and does not satisfy (8), we reduce the parameter α .

3. Descent Method with Exact Line Search

In this section we describe the descent method with an exact line search and we prove its global convergence to a solution of the problem (1).

Theorem 2. *If the assumptions (A1) – (A3) are fulfilled, then Algorithm 1 either stops at a solution of the problem (1) after a finite number of iterations, or generates a sequence $\{x^k\}$ such that any of its cluster points solves (1), or generates a sequence $\{z^i\}$, for some fixed α_k , such that any of its cluster points solves (1).*

Proof. There are three possible cases.

Case 1. The algorithm stops at x^k after a finite number of iterations. From the stopping criterion at step 1 it follows that $\varphi_{\alpha_k}(x^k) = 0$, thus x^k solves the problem (1).

Case 2. The algorithm generates an infinite sequence $\{x^k\}$. From condition at step 2b we have

$$\varphi_{\alpha_k}(x^k) \leq \frac{\alpha_k}{2(1-\eta)} \|x^k - y_{\alpha_k}(x^k)\|_G^2 \quad \forall k \in \mathbb{N}.$$

Since x^k and $y_{\alpha_k}(x^k)$ belong to X which is bounded, the norm $\|x^k - y_{\alpha_k}(x^k)\|_G^2$ is bounded above. Moreover $\lim_{k \rightarrow \infty} \alpha_k = 0$, thus

$$\lim_{k \rightarrow \infty} \varphi_{\alpha_k}(x^k) = 0. \tag{9}$$

The sequence $\{x^k\}$ has cluster points because it is bounded. Let x^* be any cluster point of $\{x^k\}$ and x^{k_s} a subsequence converging to x^* . From the definition of φ_{α_k} it follows that for each $y \in X$ we have

$$\varphi_{\alpha_{k_s}}(x^{k_s}) \geq \langle F(x^{k_s}), x^{k_s} - y \rangle - \frac{\alpha_{k_s}}{2} \|x^{k_s} - y\|_G^2 + f(x^{k_s}) - f(y) \quad \forall s \in \mathbb{N}.$$

Taking into account the continuity of F and f , $\lim_{k \rightarrow \infty} \alpha_k = 0$, and (9), then passing to the limit we obtain

$$0 \geq \langle F(x^*), x^* - y \rangle + f(x^*) - f(y).$$

Since y is arbitrary, we have proved that x^* solves the problem (1).

Case 3. The algorithm generates an infinite sequence $\{z^i\}$ for a fixed $\alpha_k = \alpha$. Let z^* be any cluster point of $\{z^i\}$ and z^{i_s} a subsequence converging to z^* . Assume by contradiction that z^* does not solve (1), thus $\varphi_\alpha(z^*) > 0$. Moreover, for all $s \in \mathbb{N}$ we have:

$$-\varphi_\alpha(z^{i_s}) + \frac{\alpha}{2} \|z^{i_s} - y_\alpha(z^{i_s})\|_G^2 < -\eta \varphi_\alpha(z^{i_s}).$$

Hence passing to the limit, since φ_α and y_α are continuous (see [4]), we obtain

$$-\varphi_\alpha(z^*) + \frac{\alpha}{2} \|z^* - y_\alpha(z^*)\|_G^2 \leq -\eta \varphi_\alpha(z^*) < 0.$$

Using Theorem 1 we have:

$$\varphi'_\alpha(z^*; y_\alpha(z^*) - z^*) < 0,$$

thus $d^* = y_\alpha(z^*) - z^*$ is a descent direction for φ_α at z^* and

$$\min_{t \in [0,1]} \varphi_\alpha(z^* + t d^*) < \varphi_\alpha(z^*). \quad (10)$$

On the other hand the sequence $\{\varphi_\alpha(z^i)\}$ is monotone decreasing and from the step length rule it follows that for each $t \in [0, 1]$ we have:

$$\varphi_\alpha(z^{i_s+1}) \leq \varphi_\alpha(z^{i_s} + t(y_\alpha(z^{i_s}) - z^{i_s})), \quad \forall s \in \mathbb{N}.$$

Passing to the limit we obtain:

$$\varphi_\alpha(z^*) \leq \varphi_\alpha(z^* + t(y_\alpha(z^*) - z^*)) \quad \forall t \in [0, 1],$$

that is

$$\varphi_\alpha(z^*) = \min_{t \in [0,1]} \varphi_\alpha(z^* + t d^*)$$

which is impossible because it contradicts (10). Thus z^* is a solution of the problem (1). \square

4. Descent Method with Armijo-Type Line Search

In this section we describe the descent method with an Armijo-type line search and we prove it is globally convergent to a solution of (1).

Theorem 3. *If the assumptions (A1)–(A3) are fulfilled, then Algorithm 2 either stops at a solution of the problem (1) after a finite number of iterations, or generates a bounded sequence $\{x^k\}$ such that any of its cluster points solves (1), or generates a bounded sequence $\{z^i\}$, for some fixed α_k , such that any of its cluster points solves (1).*

Proof. First, we show that the algorithm is well defined, i.e. that the line search procedure is always finite. Assume, by contradiction, that there are $i, k \geq 0$ such that the inequality

$$\varphi_{\alpha_k}(z^i + \gamma^m d^i) - \varphi_{\alpha_k}(z^i) > -\beta \gamma^m \varphi_{\alpha_k}(z^i),$$

holds for all $m \in \mathbb{N}$. Then we have:

$$\varphi'_{\alpha_k}(z^i; d^i) = \lim_{m \rightarrow +\infty} \frac{\varphi_{\alpha_k}(z^i + \gamma^m d^i) - \varphi_{\alpha_k}(z^i)}{\gamma^m} \geq -\beta \varphi_{\alpha_k}(z^i).$$

Algorithm 2

0. (Initial step)

Let G be a symmetric positive definite matrix, $\eta, \gamma \in (0, 1)$, and $\beta \in (0, \eta)$.

Let $\{\alpha_k\}$ be a sequence strictly decreasing to 0.

Choose any $x^0 \in X$ and set $k = 0$.

1. (Stopping criterion)

If $\varphi_{\alpha_k}(x^k) = 0$

then STOP,

else set $k = k + 1$.

2. (Minimization of φ_{α_k})

2a. (Initialization)

Set $i = 0$ and $z^0 = x^{k-1}$.

2b. **If** $-\varphi_{\alpha_k}(z^i) + \frac{\alpha_k}{2} \|z^i - y_{\alpha_k}(z^i)\|_G^2 < -\eta \varphi_{\alpha_k}(z^i)$

then (line search)

set $d^i = y_{\alpha_k}(z^i) - z^i$

compute the smallest nonnegative integer m such that:

$$\varphi_{\alpha_k}(z^i + \gamma^m d^i) - \varphi_{\alpha_k}(z^i) \leq -\beta \gamma^m \varphi_{\alpha_k}(z^i)$$

set $t_i = \gamma^m$,

else (update of x^k)

set $x^k = z^i$ and return to step 1.

2c. (Update of z^i)

Set $z^{i+1} = z^i + t_i d^i$, $i = i + 1$, and return to step 2b.

Combining (4) and step 2b, we get:

$$\varphi'_{\alpha_k}(z^i; d^i) \leq -\varphi_{\alpha_k}(z^i) + \frac{\alpha_k}{2} \|d^i\|_G^2 < -\eta \varphi_{\alpha_k}(z^i),$$

therefore

$$(\eta - \beta) \varphi_{\alpha_k}(z^i) < 0,$$

which is impossible because $\eta > \beta$ and $\varphi_{\alpha_k}(z^i) \geq 0$. So the line search procedure is always finite.

There are three possible cases.

Case 1. The algorithm stops at x^k after a finite number of iterations. From the stopping criterion it follows that x^k solves (1).

Case 2. The algorithm generates an infinite sequence $\{x^k\}$. As in the Case 2 of Theorem 2, it can be proved that any cluster point of $\{x^k\}$ solves (1).

Case 3. The algorithm generates an infinite sequence $\{z^i\}$ for a fixed $\alpha_k = \alpha$. Let us consider two possible subcases: either $\limsup_{i \rightarrow \infty} t_i > 0$, or $\limsup_{i \rightarrow \infty} t_i = 0$.

Subcase 3a. If $\limsup_{i \rightarrow \infty} t_i > 0$, then there exists $t^* > 0$ and a subsequence $\{t_{i_s}\}$ such that $t_{i_s} \geq t^* > 0$ for all $s \in \mathbb{N}$. Since the sequence $\{z^i\}$ is infinite, we have:

$$\varphi_\alpha(z^{i_s}) - \varphi_\alpha(z^{i_s+1}) \geq \beta t_{i_s} \varphi_\alpha(z^{i_s}) \geq \beta t^* \varphi_\alpha(z^{i_s}) \geq 0. \tag{11}$$

The sequence $\{\varphi_\alpha(z^i)\}$ is monotone decreasing and bounded below, hence

$$\lim_{i \rightarrow \infty} [\varphi_\alpha(z^i) - \varphi_\alpha(z^{i+1})] = 0,$$

and in particular

$$\lim_{s \rightarrow \infty} [\varphi_\alpha(z^{i_s}) - \varphi_\alpha(z^{i_s+1})] = 0. \tag{12}$$

Using (11) and (12), we obtain $\lim_{s \rightarrow \infty} \varphi_\alpha(z^{i_s}) = 0$ and thus $\lim_{i \rightarrow \infty} \varphi_\alpha(z^i) = 0$. If z^* is any cluster point of $\{z^i\}$, then from the continuity of φ_α it follows that $\lim_{i \rightarrow \infty} \varphi_\alpha(z^i) = \varphi_\alpha(z^*)$, hence $\varphi_\alpha(z^*) = 0$, i.e. z^* is a solution of the problem (1).

Subcase 3b. If $\limsup_{i \rightarrow \infty} t_i = 0$, then $\lim_{i \rightarrow \infty} t_i = 0$. From the step length rule it follows that for all $i \in \mathbb{N}$,

$$\varphi_\alpha(z^i + \gamma^{-1} t_i d^i) - \varphi_\alpha(z^i) > -\beta \gamma^{-1} t_i \varphi_\alpha(z^i).$$

By the mean value theorem we have

$$\varphi_\alpha(z^i + \gamma^{-1} t_i d^i) - \varphi_\alpha(z^i) = \langle \xi^i, \gamma^{-1} t_i d^i \rangle,$$

where $\xi^i \in \partial \varphi_\alpha(z^i + \theta_i \gamma^{-1} t_i d^i)$ for some $\theta_i \in (0, 1)$. We set $w^i = z^i + \theta_i \gamma^{-1} t_i d^i$. From [10, Proposition 4.19] it follows that

$$\xi^i = F(w^i) - (\nabla F(w^i)^\top - \alpha G)(y_\alpha(w^i) - w^i) + g^i,$$

for some $g^i \in \partial f(w^i)$. Therefore, for all $i \in \mathbb{N}$, we have:

$$\langle F(w^i) - (\nabla F(w^i)^\top - \alpha G)(y_\alpha(w^i) - w^i), d^i \rangle + \langle g^i, d^i \rangle > -\beta \varphi_\alpha(z^i).$$

The sequences $\{z^i\}$ and $\{d^i\}$ are bounded, thus also $\{g^i\}$ is bounded. Let z^* be any cluster point of $\{z^i\}$. Since $\lim_{i \rightarrow \infty} t_i = 0$ and the set-valued map ∂f is closed, passing to the limit and taking a subsequence if necessary, we get:

$$\langle F(z^*) - (\nabla F(z^*)^\top - \alpha G)(y_\alpha(z^*) - z^*), d^* \rangle + \langle g^*, d^* \rangle \geq -\beta \varphi_\alpha(z^*), \tag{13}$$

where $d^* = y_\alpha(z^*) - z^*$ and $g^* \in \partial f(z^*)$. Since

$$f'(z^*; d^*) = \max_{g \in \partial f(z^*)} \langle g, d^* \rangle, \tag{14}$$

from [10, Proposition 4.19], (13) and (14) it follows that:

$$\begin{aligned} \varphi'_\alpha(z^*; d^*) &= \langle F(z^*) - (\nabla F(z^*)^\top - \alpha G)(y_\alpha(z^*) - z^*), d^* \rangle + f'(z^*; d^*) \\ &\geq \langle F(z^*) - (\nabla F(z^*)^\top - \alpha G)(y_\alpha(z^*) - z^*), d^* \rangle + \langle g^*, d^* \rangle \geq -\beta \varphi_\alpha(z^*). \end{aligned} \tag{15}$$

Moreover, for all $i \in \mathbb{N}$, we have:

$$-\varphi_\alpha(z^i) + \frac{\alpha}{2} \|z^i - y_\alpha(z^i)\|_G^2 < -\eta \varphi_\alpha(z^i),$$

hence passing to the limit and taking a subsequence if necessary, and using Theorem 1 we obtain:

$$\varphi'_\alpha(z^*; d^*) \leq -\varphi_\alpha(z^*) + \frac{\alpha}{2} \|d^*\|_G^2 \leq -\eta \varphi_\alpha(z^*). \tag{16}$$

From (15) and (16) we get

$$(\eta - \beta) \varphi_\alpha(z^*) \leq 0.$$

Since $\eta > \beta$ and $\varphi_\alpha(z^*) \geq 0$, it follows that $\varphi_\alpha(z^*) = 0$. i.e. z^* solves the problem (1). □

Remark 4. In Algorithms 1 and 2 the sequence $\{\alpha_k\}$ can be chosen adaptively, for example (see also [11]) such as:

$$\alpha_k = \begin{cases} \alpha_{k-1} & \text{if } \varphi_{\alpha_{k-1}}(x^{k-1}) \leq \nu_{k-1}, \\ \mu \alpha_{k-1} & \text{otherwise,} \end{cases} \tag{17}$$

where $0 < \mu < 1$ and $\{\nu_k\}$ is a sequence decreasing to 0. Indeed, if the algorithm generates an infinite sequence $\{x^k\}$ with $\{\alpha_k\}$ chosen by (17), then either $\lim_{k \rightarrow \infty} \alpha_k = 0$, which can be treated as in the Case 2 of Theorem 2 or Theorem 3, or one has

$$\alpha_k = \bar{\alpha} \quad \text{and} \quad \varphi_{\bar{\alpha}}(x^k) \leq \nu_k \quad \forall k > \bar{k},$$

hence $\lim_{k \rightarrow \infty} \varphi_{\bar{\alpha}}(x^k) = 0$. Then for each cluster point x^* of $\{x^k\}$ we have $\varphi_{\bar{\alpha}}(x^*) = 0$, that is x^* solves the problem (1).

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